Windows Defender Under The Microscope:

A Reverse Engineer's Perspective

Alexei Bulazel
@0xAlexei

Virus Bulletin 2018

About Me

- AV industry outsider working on AV RE for a long time
- Security researcher at ForAllSecure
- RPI / RPISEC alumnus
  - Co-taught the famous RPISEC “Modern Binary Exploitation” class (`https://github.com/rpisec/mbe`)
- First time at Virus Bulletin

This is my personal research, any views and opinions expressed are my own, not those of any employer

@0xAlexei

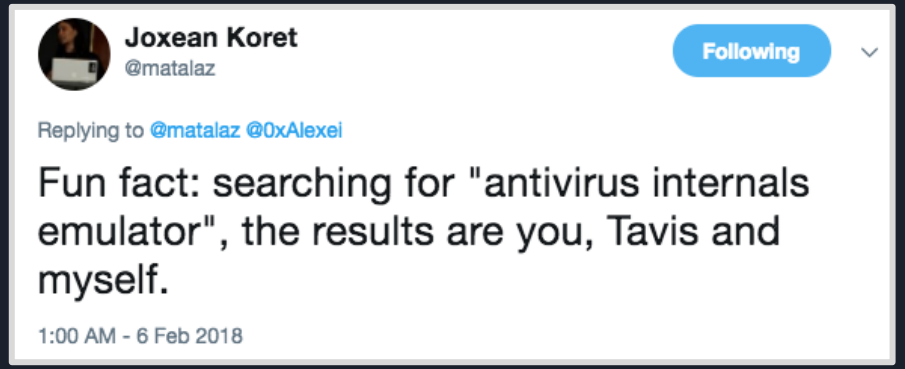
RPISEC

Outline

1. Introduction
2. Tooling & Process
3. Discussion
4. Conclusion

This Talk

- Analysis of my custom tools and process after 9+ months of REing Windows Defender
  - Not going to reiterate AV knowledge that industry already knows - see released slides



- Few researchers REing AVs, fewer looking at emulators
- No disrespect to Microsoft or the AV industry - Defender is a fascinating subject of study and a beautifully architected piece of software

My Published Research

Windows Defender RE

- JS Engine @ REcon Brussels
- Windows Emulator @ REcon Montreal, Black Hat, DEFCON

“AVLeak” - AV emulator fingerprinting and evasion @ Black Hat & WOOT’16

“A Survey On Automated Dynamic Malware Analysis Evasion and Counter-Evasion” @ ROOTS’17

Reverse Engineering Windows Defender’s JavaScript Engine

Reverse Engineering Windows Defender Part II: The Windows Binary Emulator

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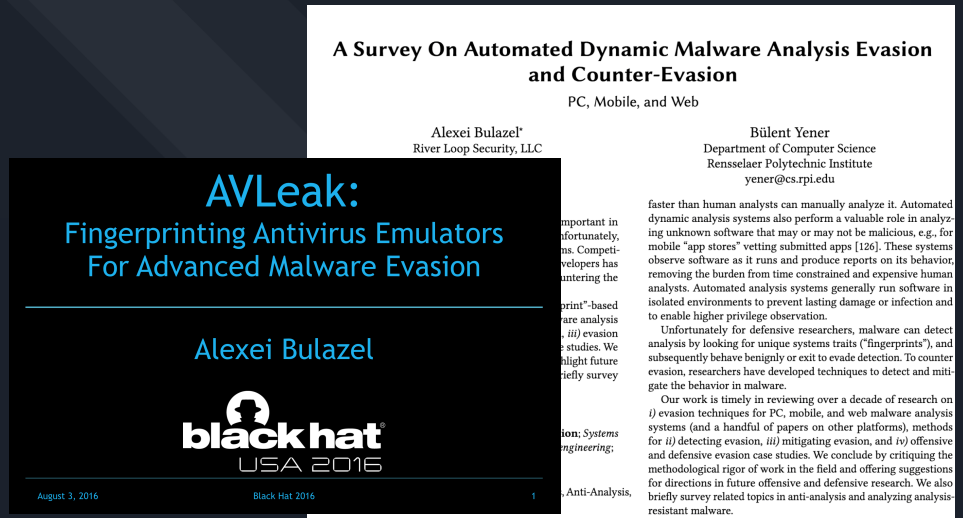
A Survey On Automated Dynamic Malware Analysis Evasion and Counter-Evasion

PC, Mobile, and Web

Alexei Bulazel*
River Loop Security, LLC

Bülent Yener
Department of Computer Science
Rensselaer Polytechnic Institute
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faster than human analysts can manually analyze it. Automated dynamic analysis systems also perform a valuable role in analyzing unknown software that may or may not be malicious, e.g., for mobile “app stores” vetting submitted apps [126]. These systems observe software as it runs and produce reports on its behavior, removing the burden from time constrained and expensive human analysts. Automated analysis systems generally run software in isolated environments to prevent lasting damage or infection and to enable higher privilege observation.

Unfortunately for defensive researchers, malware can detect analysis by looking for unique systems traits (“fingerprints”), and subsequently behave benignly or exit to evade detection. To counter evasion, researchers have developed techniques to detect and mitigate the behavior in malware.

Our work is timely in reviewing over a decade of research on *i)* evasion techniques for PC, mobile, and web malware analysis systems (and a handful of papers on other platforms), methods for *ii)* detecting evasion, *iii)* mitigating evasion, and *iv)* offensive and defensive evasion case studies. We conclude by critiquing the methodological rigor of work in the field and offering suggestions for directions in future offensive and defensive research. We also briefly survey related topics in anti-analysis and analyzing analysis-resistant malware.

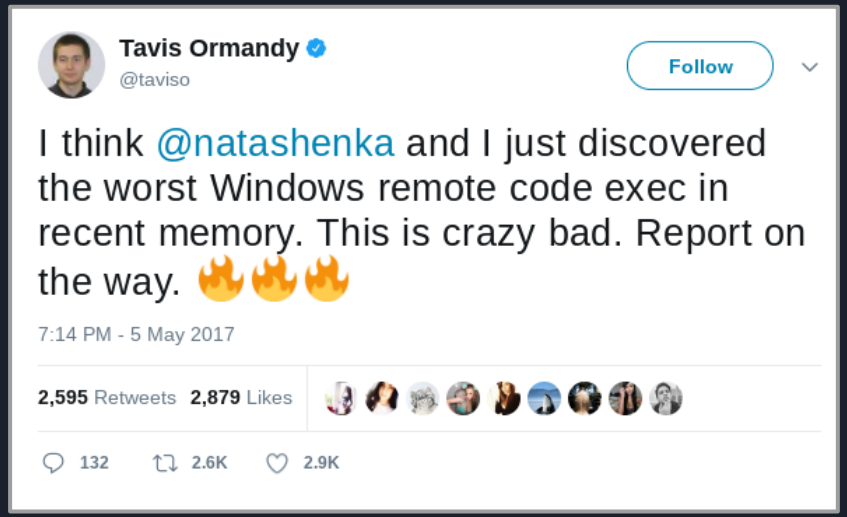
AVLeak:
Fingerprinting Antivirus Emulators For Advanced Malware Evasion

Alexei Bulazel

black hat USA 2016

August 3, 2016 Black Hat 2016 1

Motivation



- Tavis and co. at P0 dropped some awesome Defender JS engine bugs
- I had analyzed AVs before, but never Windows Defender… interest in JS engines
- So I reverse engineered Defender's JS engine for ~4 months
- I then spent another ~5 months reverse engineering the Windows binary emulator
- This was a *personal* research project - all in my free time, not for any company

Real Motivation

Spend hundreds of hours doing unpaid research, so I can fly thousands of miles in coach class to present Powerpoints in hotels around the world

Prior Art

- Lots of conference talks, whitepapers, and blogs on antivirus *evasion*, but few on RE
- Tavis Ormandy's Defender bugs from 2017
- As far as I know, there's never been a publication about *reverse engineering* the internals of an AV emulator*

*There are plenty on black box AV evasion though. AV industry companies have occasionally presented on the design of their emulators at conferences such as Virus Bulletin.

MsMpEng: Multiple problems handling ntdll!NtControlChannel commands

Project Member Reported by taviso@google.com, May 12 2017

MsMpEng includes a full system x86 emulator that is used to execute any untrusted files that lo
runs as NT AUTHORITY\SYSTEM and isn't sandboxed.

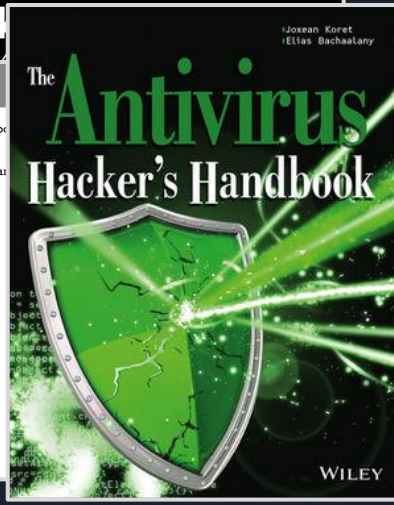
Browsing the list of win32 APIs that the emulator supports, I noticed ntdll!NtControlChannel, a
emulated code to control the emulator.

You can simply create an import library like this and then call it from emulated code:

```
$ cat ntdll.def
LIBRARY ntdll.dll
EXPORTS
    NtControlChannel
$ lib /def:ntdll.def /machine:x86 /out:ntdll.lib /nologo
   Creating library ntdll.lib and object ntdll.exp
$ cat intoverflow.c
#include <windows.h>
#include <stdint.h>
#include <stdlib.h>
#include <limits.h>

#pragma pack(1)

struct {
    uint64_t start_va;
    uint32_t size;
```



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Reconnaissance - Patent Search

(12) **United States Patent**
Gheorghescu et al.

(10) **Patent No.: US 7,636,856 B2**
(45) **Date of Patent: Dec. 22, 2009**

(54) **PROACTIVE COMPUTER MALWARE PROTECTION THROUGH DYNAMIC TRANSLATION**

(75) Inventors: **Gheorghe Marius Gheorghescu**, Redmond, WA (US); **Adrian M Marinescu**, Sammamish, WA (US); **Adrian E Stepan**, Redmond, WA (US)

(73) Assignee: **Microsoft Corporation**, Redmond, WA

6,330,691	B1 *	12/2001	Buzbee et al.	714/35
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6,704,925	B1 *	3/2004	Bugnion	717/138
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2003/0101381	A1 *	5/2003	Mateev et al.	714/38
2005/0005153	A1 *	1/2005	Das et al.	713/200

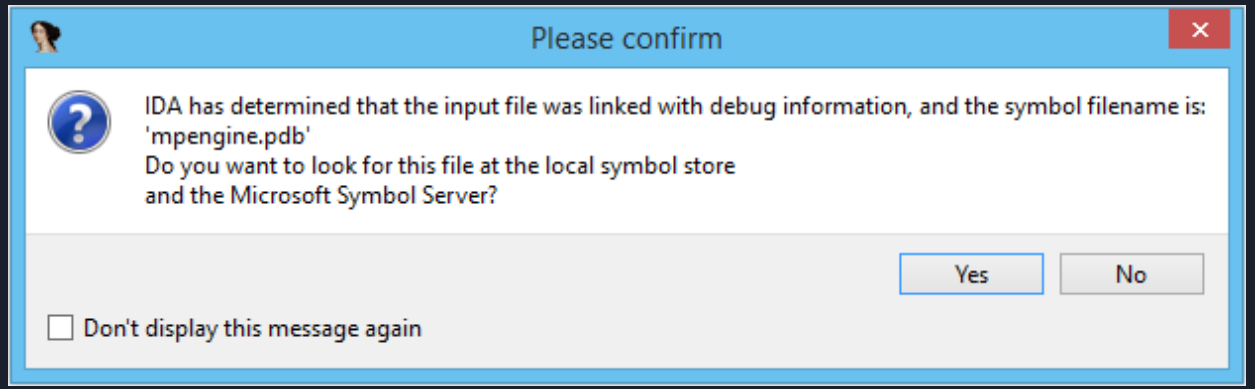
OTHER PUBLICATIONS

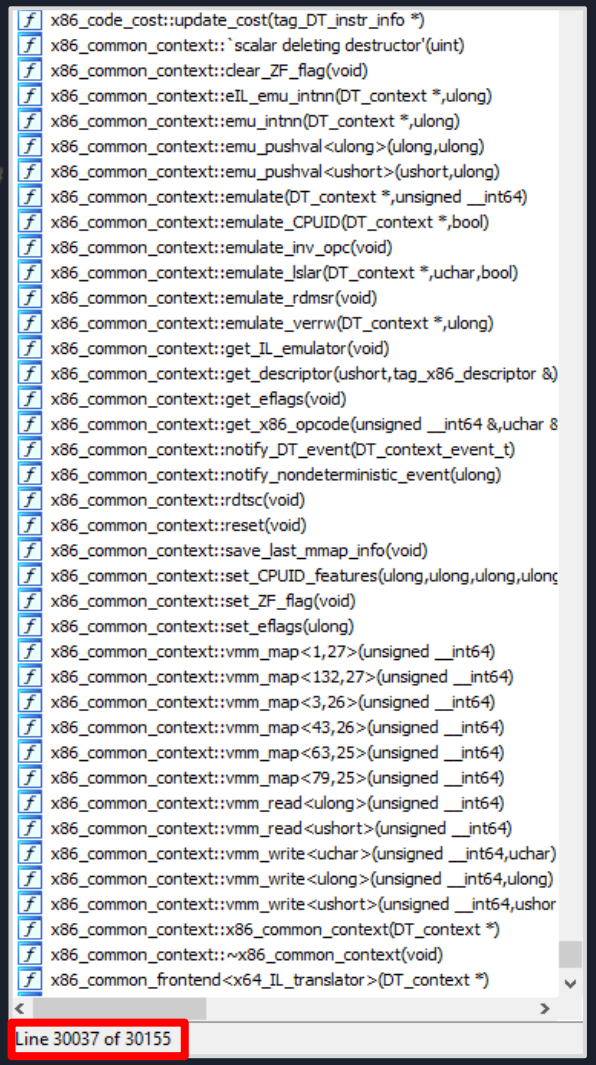
Cifuentes, Cristina, "Reverse Compilation Techniques," Jul. 1994

"The present invention includes a system and method for translating potential malware devices into safe program code. The potential malware is translated from any one of a number of different types of source languages, including, but not limited to, native CPU program code, platform independent .NET byte code, scripting program code, and the like. Then the translated program code is compiled into program code that may be understood and executed by the native CPU…"

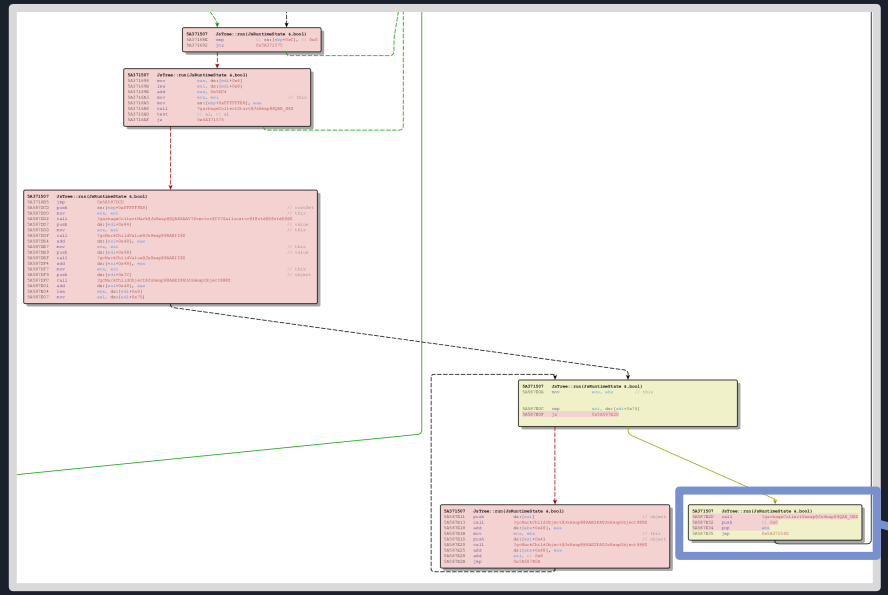
Static Analysis

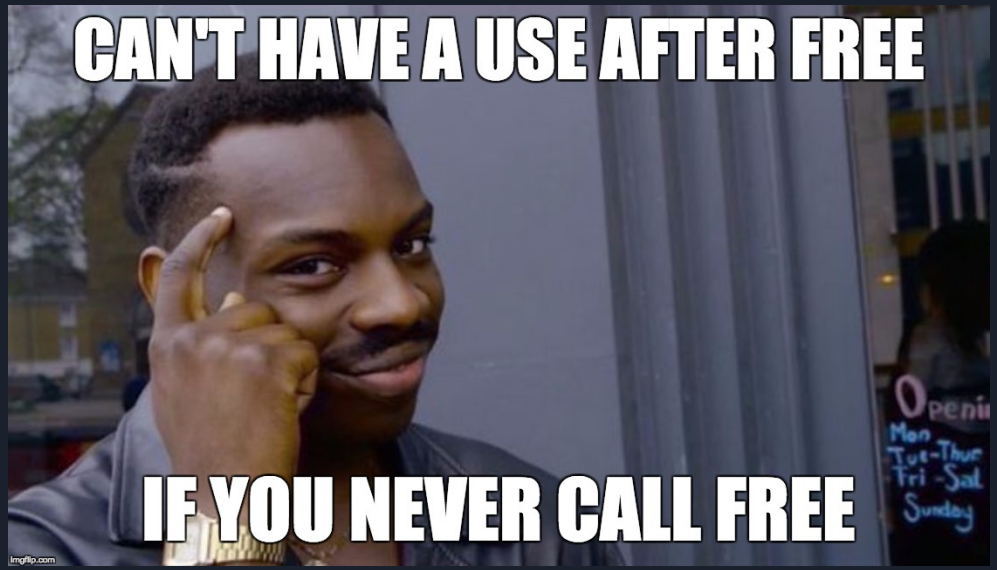
- ~12 MB DLL
- ~30,000 functions
- IDA Pro
  - Patch analysis with BinDiff
- Microsoft publishes PDBs

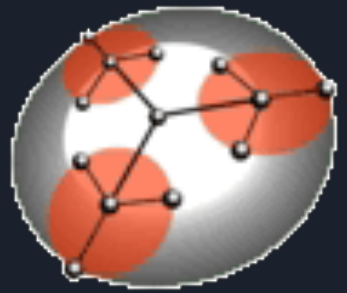




BinDiffing





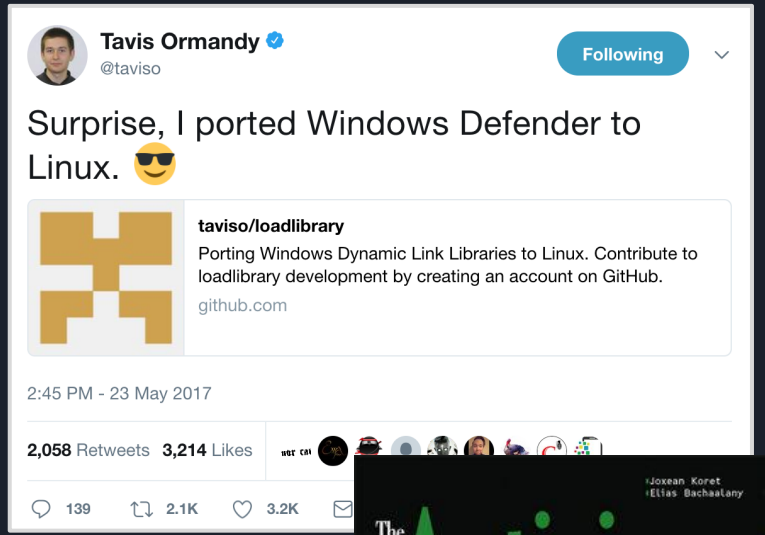


```
5A371507    JsTree::run(JsRuntimeState &,bool)
5A587E2D    call        ?garbageCollectSweep@JsHeap@@QAE_NXZ
5A587E32    push        b1 0x0
5A587E34    pop         ebx
5A587E35    jmp         0x5A37156D
```

Dynamic Analysis & Loader

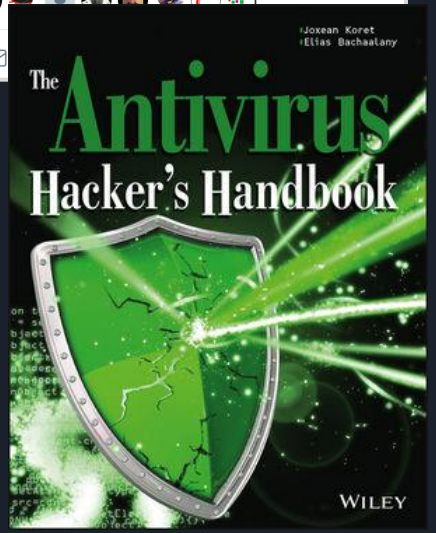
AV-Specific Challenges:

- Protected Process
  - Cannot debug, even as local admin
- Introspection
- Scanning on demand
- Code reachability may be configuration / heuristics dependent



Example: MPEngine Lockdown

- “Protected Processes” - Windows programs that you cannot debug with a usermode debugger, even if you have all privileges
- Attackers can load a signed vulnerable driver, run an exploit, get execution & deprotect the process - so … why?



“Repeated vs. single-round games in security”
Halvar Flake, BSides Zurich Keynote

Dynamic Analysis & Loader

AV-Specific Challenges:

- Protected Process
  - Cannot debug, even as local admin
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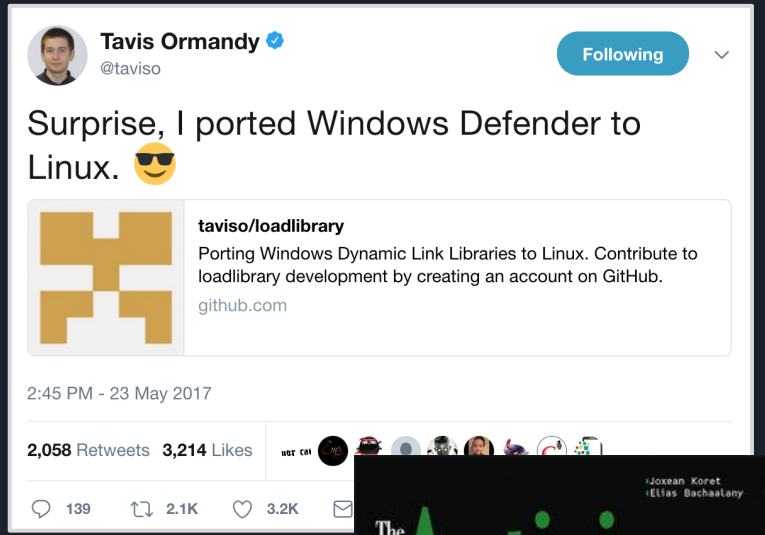
- Scanning on demand
- Code reachability may be configuration / heuristics dependent

Solution:

Custom loaders for AV binaries



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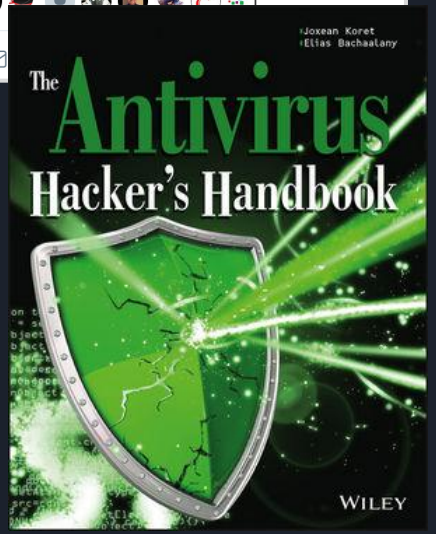

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JS REPL Shell

```
$ ./JsShell.exe
CONSTRUCTOR_CALL:       6EA109AE
DESTRUCTOR:             6EA21830
CONSTRUCTOR:            6EA21ACA
EVAL:                   6EA10875

mpscript>  (function (){for(var i = 0; i < 3; i++){print(i + ": Hello from insid
e MpEngine.dll")}})()
print(): 0: Hello from inside MpEngine.dll
print(): 1: Hello from inside MpEngine.dll
print(): 2: Hello from inside MpEngine.dll
print(): undefined
Log():              <NA>:   0: execution took 239 ticks
Log():              <NA>:   0: final memory used 9KB
Log():              <NA>:   0: total of 0 GCs performed

Ended. Result code: 0
mpscript> _
```

Based off a shell released on Twitter by `@TheWackOlian`, developed with Rolf Rolles

JS Loader and Shell

```
JsRuntimeState::triggerEvent(jsState, 0, "print", strCstr, strCstr_4, v8, v8)
```

- Use `LoadLibrary` on Windows
  - WinDbg works natively
- Patch constructor for `JsRuntimeState::JsRuntimeState()`
  - Provide a VTable implementing analysis callbacks
  - Print to `stdout` on "`print`" events
  - Log other events
- Directly call to start scan:

```
JavaScriptInterpreter::eval(
    const char *input,
    unsigned int inputSize,
    JavaScriptInterpreter::Params *params)
```

```
mov     esi, [ebp+toStringTree.baseclass_0.vfptr]
push    ecx             ; monitor
lea     ecx, [ebp+jsState] ; this
push    dword ptr [esi+20h] ; domWrapper
push    dword ptr [esi+14h] ; regexpLimit
push    dword ptr [esi+18h] ; gcLimit
push    dword ptr [esi+10h] ; memLimit
push    dword ptr [esi+0Ch] ; exeLimit
call    ??0JsRuntimeState@@QAE@IIIIPAUHtmlDocumentProvider@@PAUJsEvaluationMonitor@@@Z
mov     byte ptr [ebp+var_4], 3
mov     ecx, [esi]
mov     al, cl
shr     al, 1
and     cl, 1
and     al, 1
mov     dl, cl          ; addBrowserRt
push    eax             ; addDomRt
lea     ecx, [ebp+jsState] ; jsState
call    ?declareGlobalProperties@@YA_NAAUJsRuntimeState@@_N1@Z ; declareGlobalPropertie
pop     ecx
test    al, al
jz      loc_5A5838CC
```

JS Loader and Shell

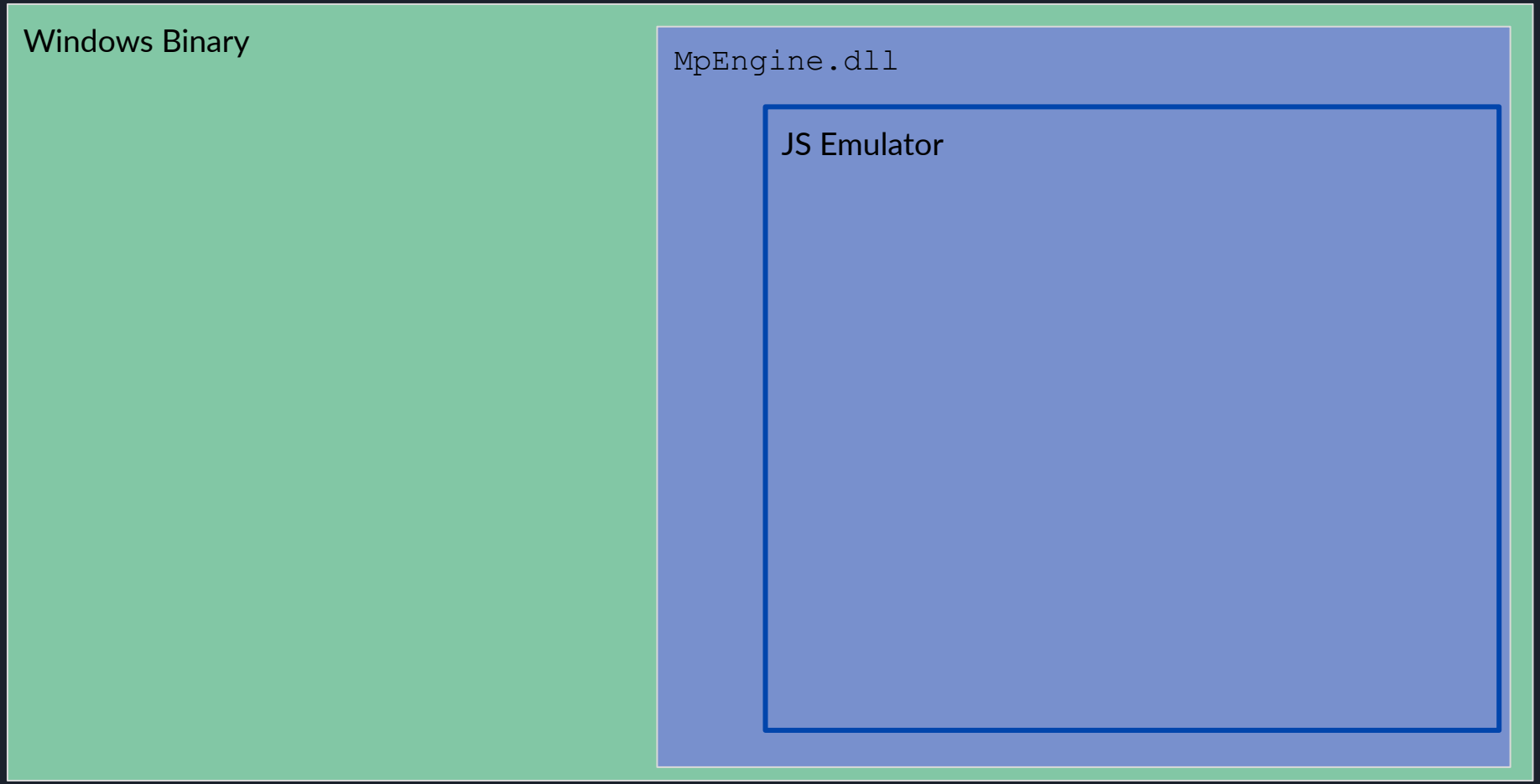
Windows Binary

JS Loader and Shell

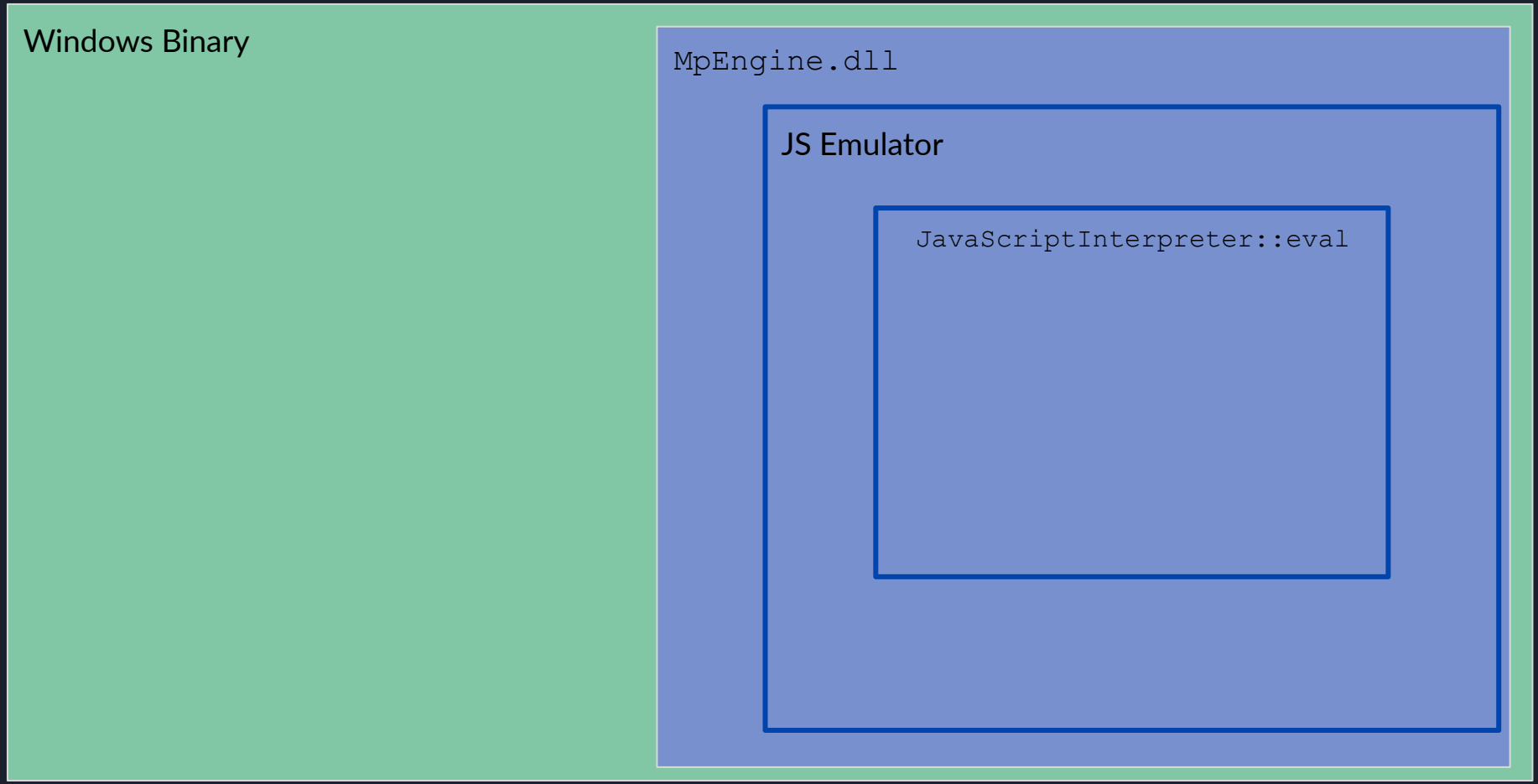
Windows Binary

MpEngine.dll

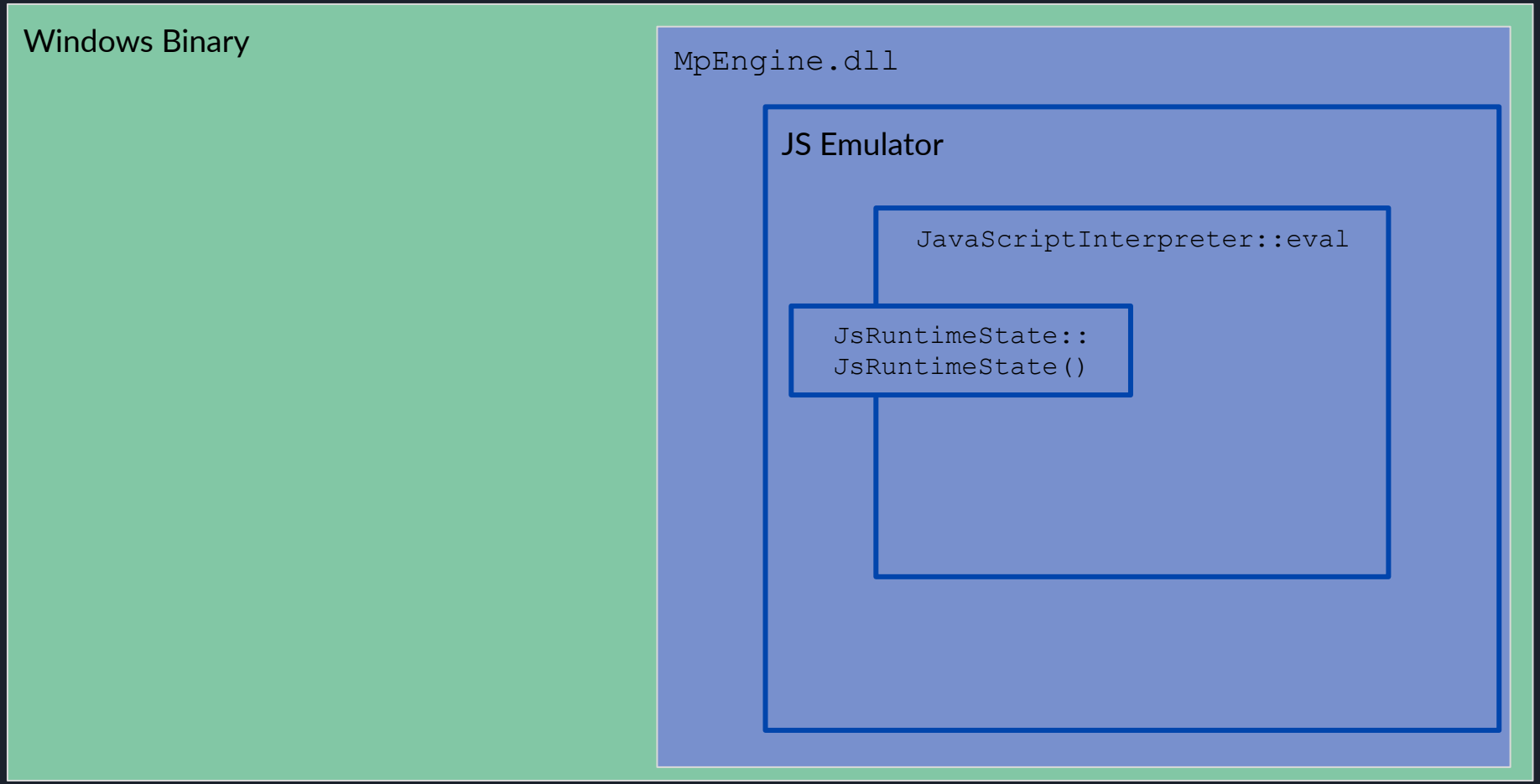
JS Loader and Shell



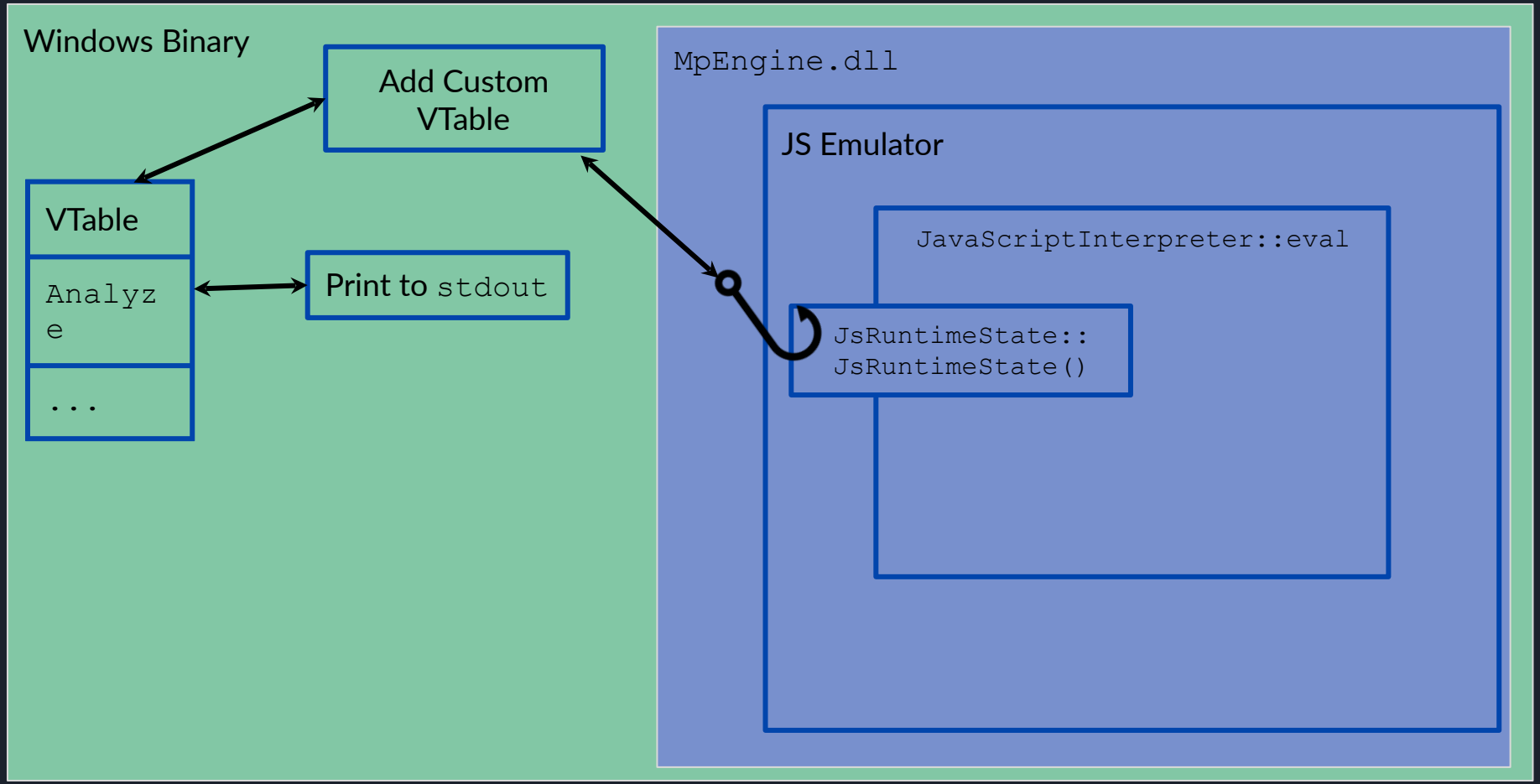
JS Loader and Shell



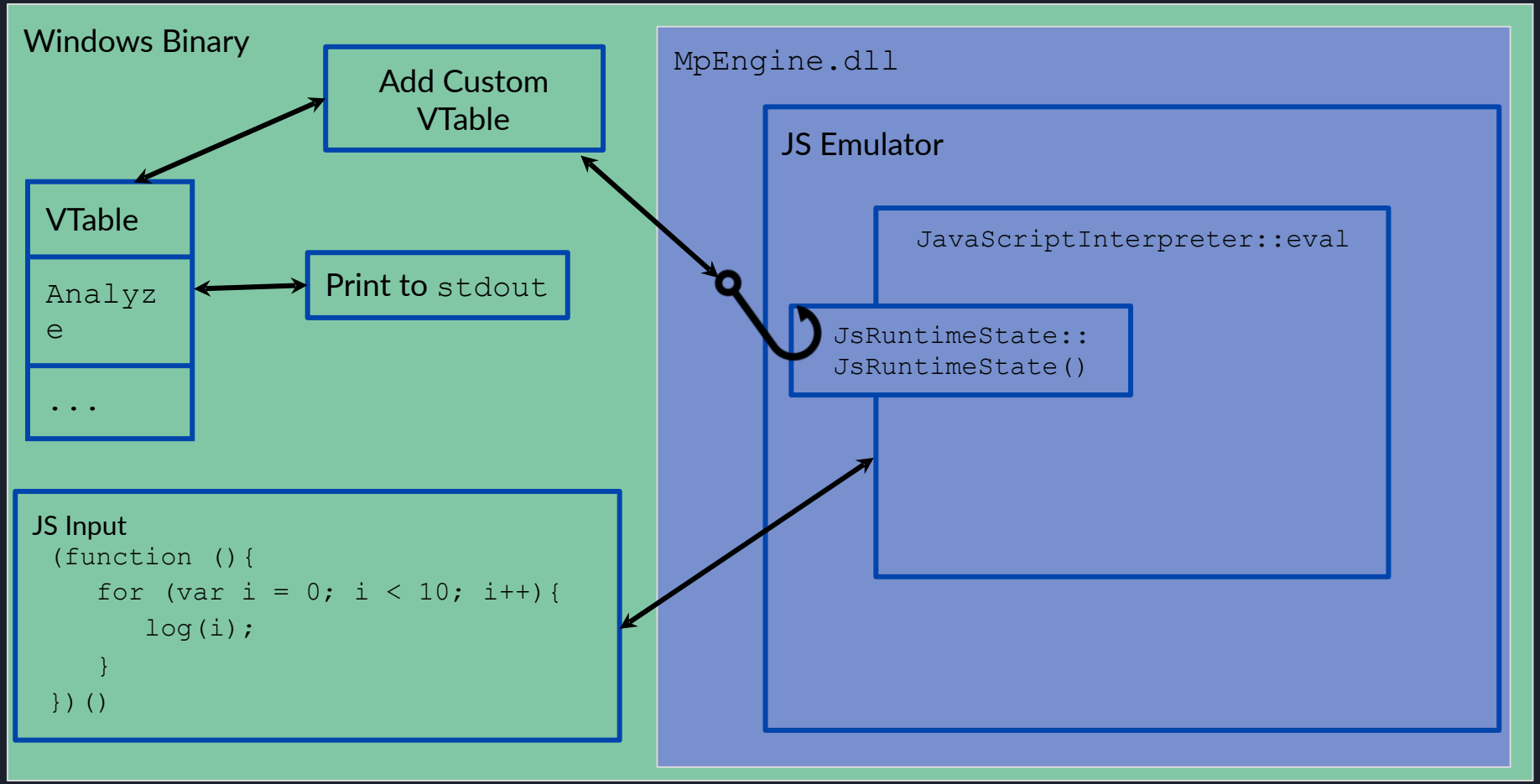
JS Loader and Shell



JS Loader and Shell



JS Loader and Shell



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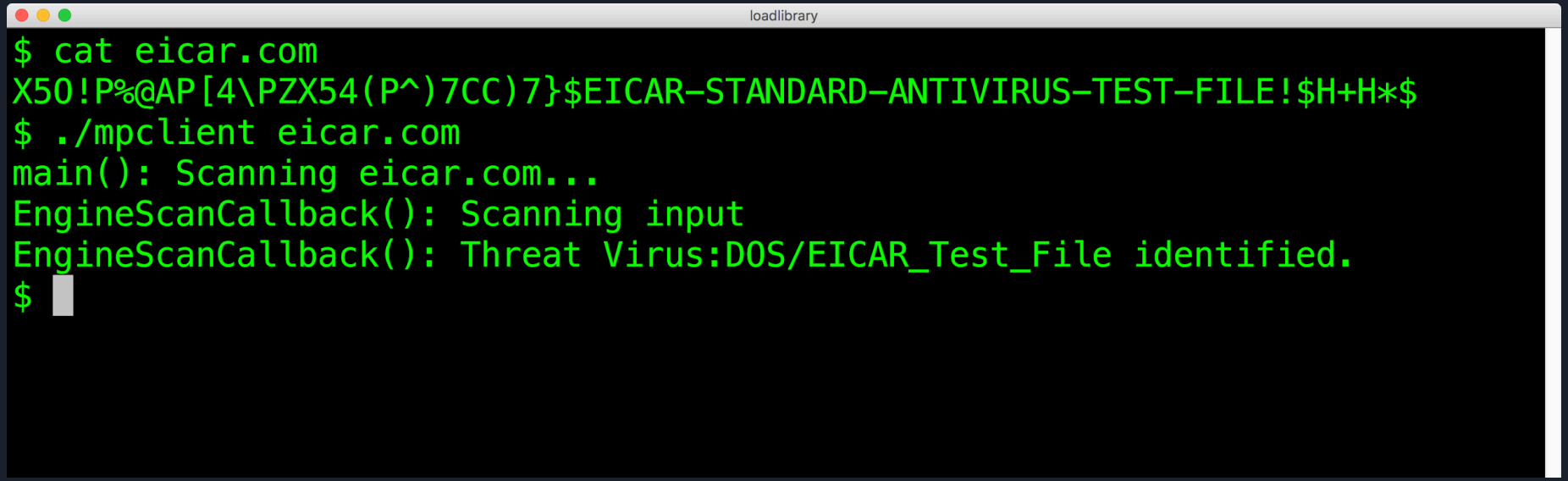
mpclient **Shell** git.io/fbp0X



Tavis Ormandy's open source tool

mpclient git.io/fbp0X

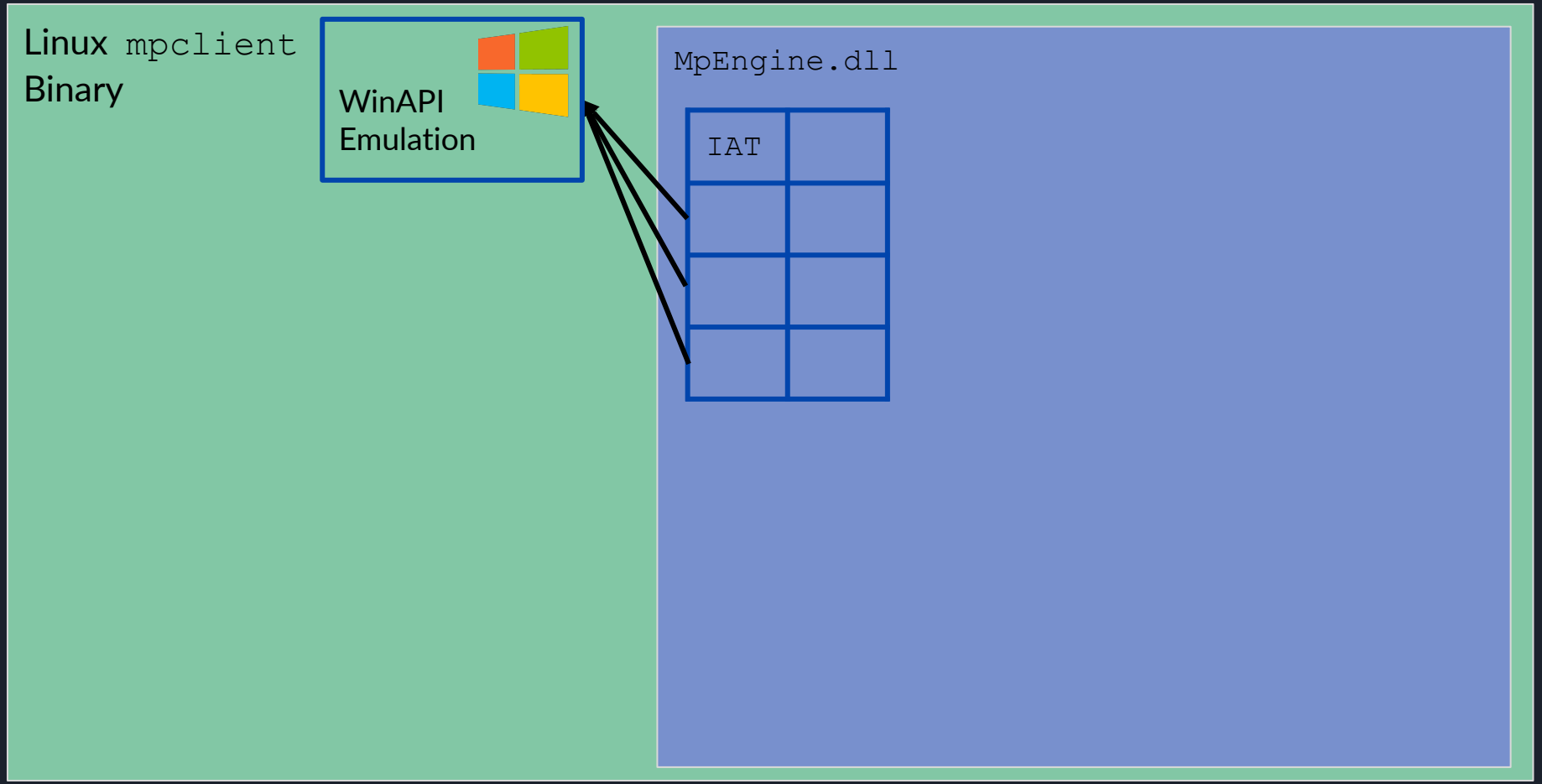
Linux `mpclient`
Binary

mpclient git.io/fbp0X

Linux `mpclient`
Binary

`MpEngine.dll`

mpclient git.io/fbp0X



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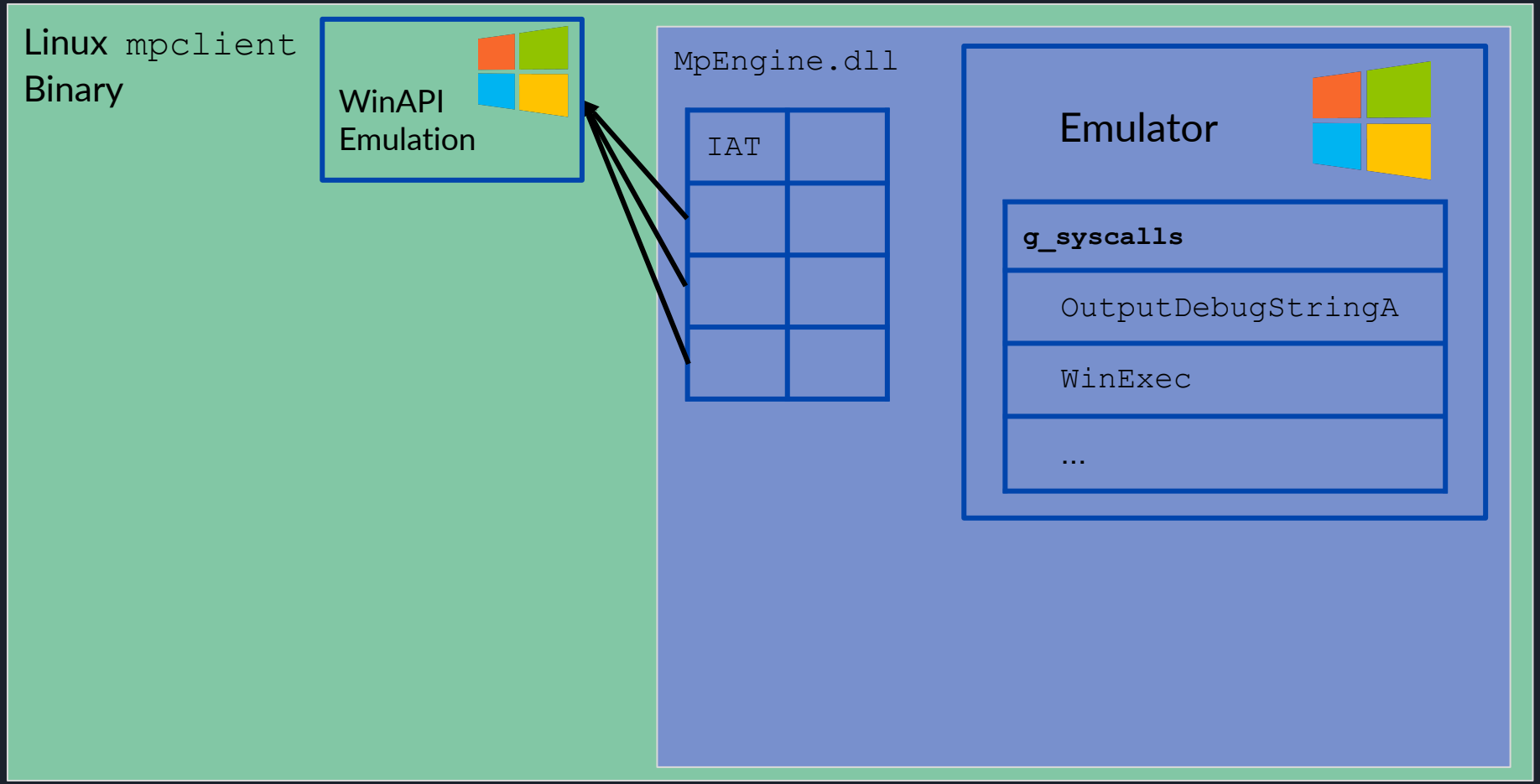

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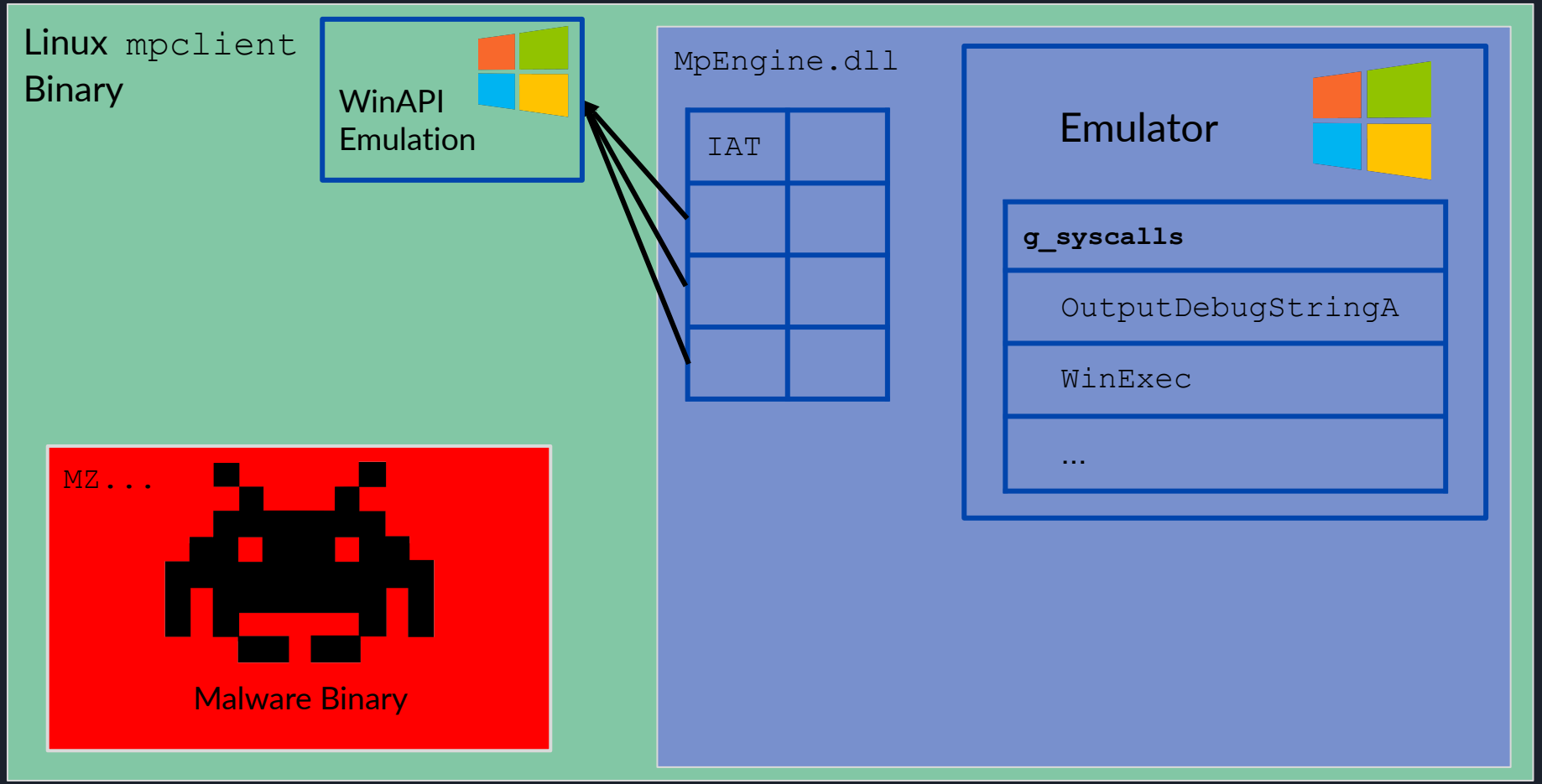
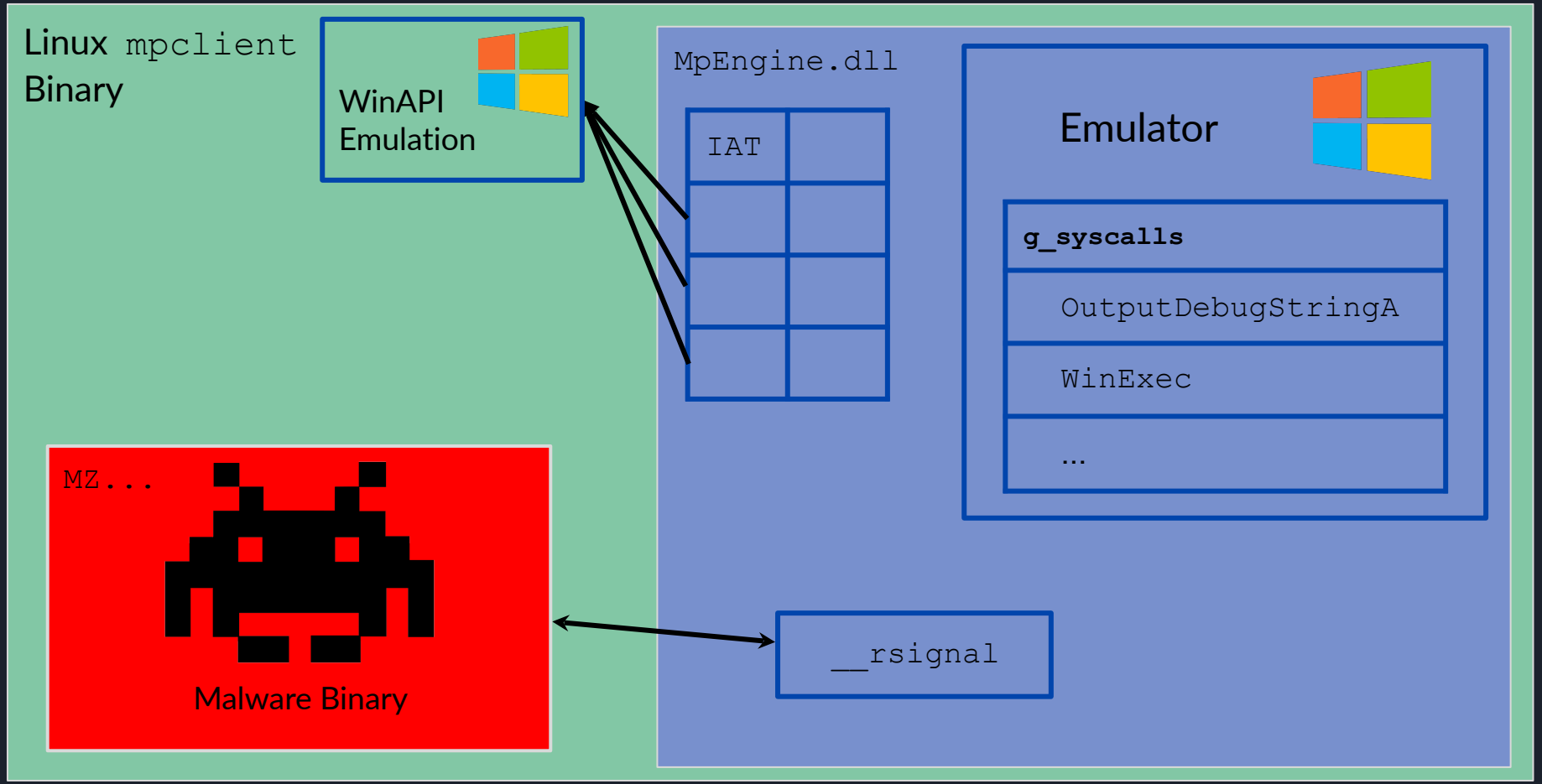


mpclient git.io/fbp0X



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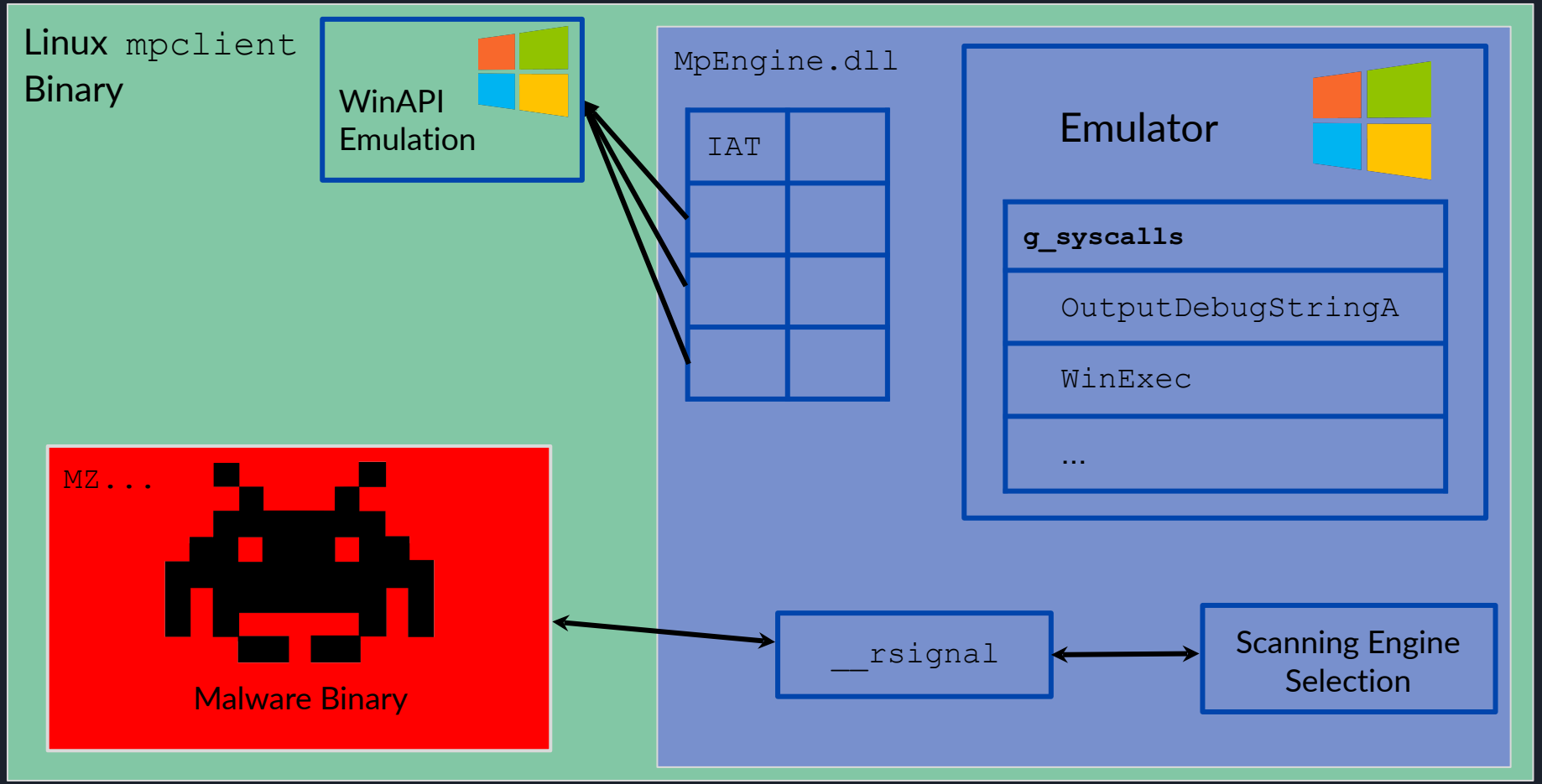
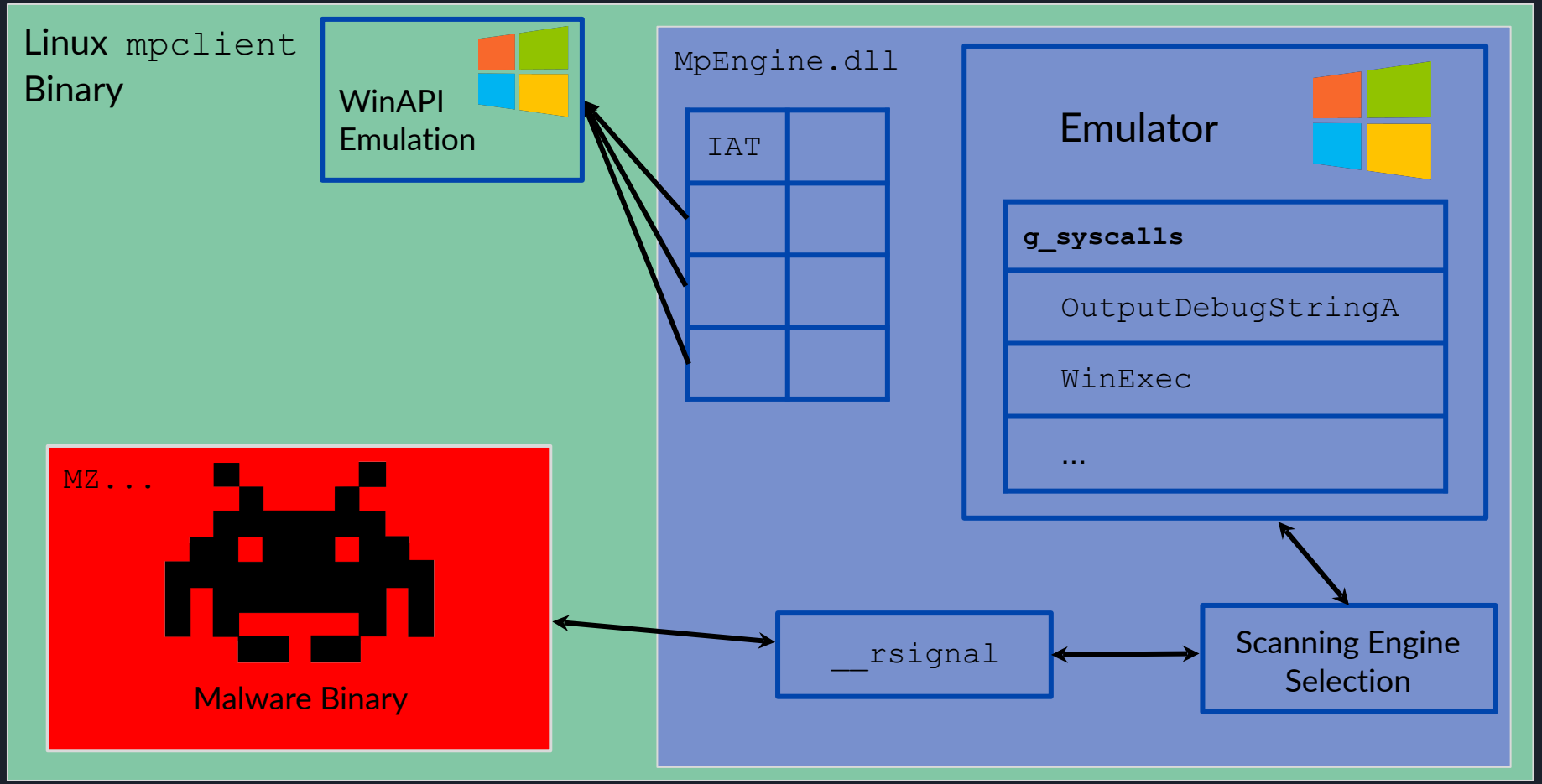
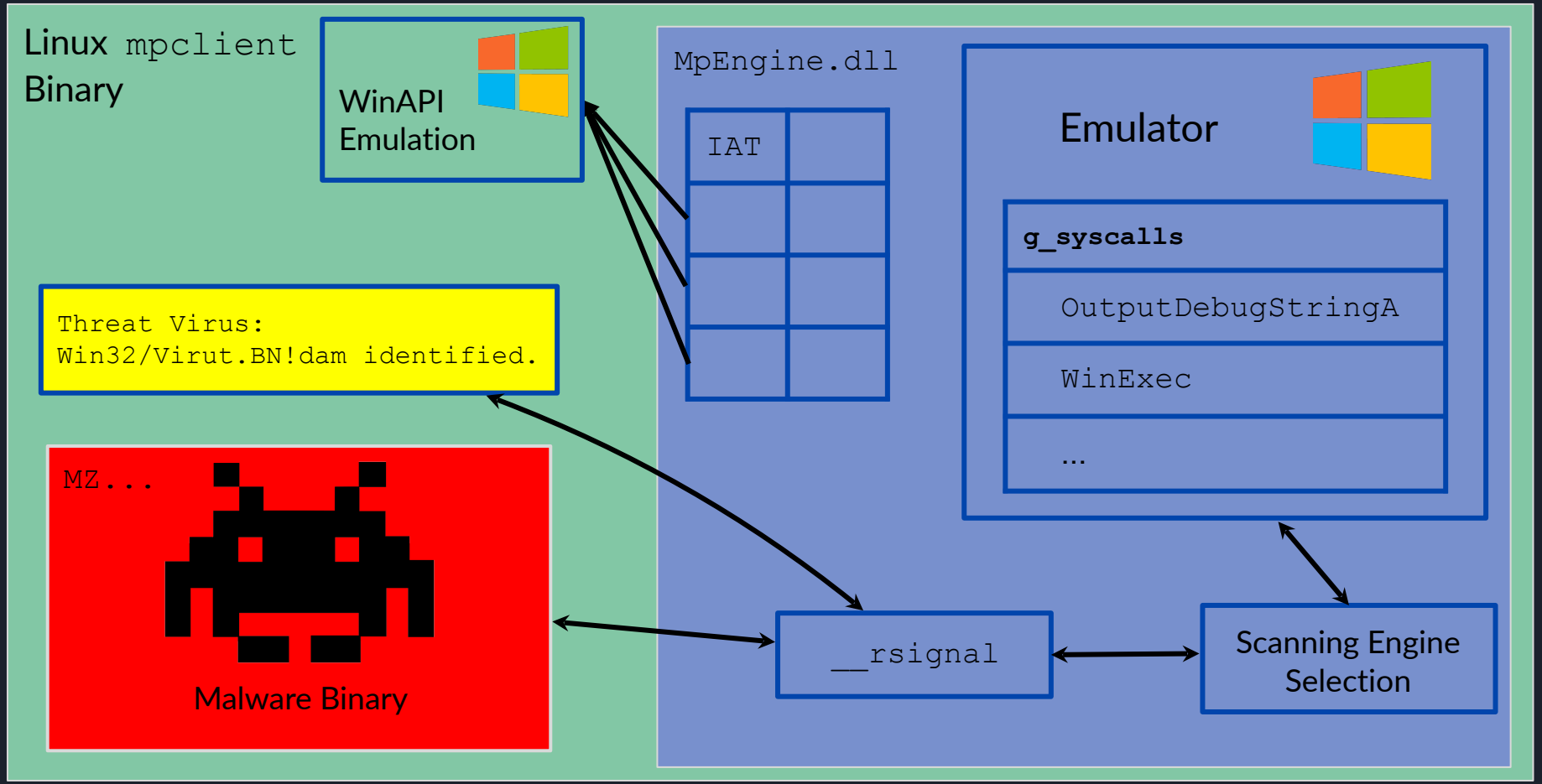


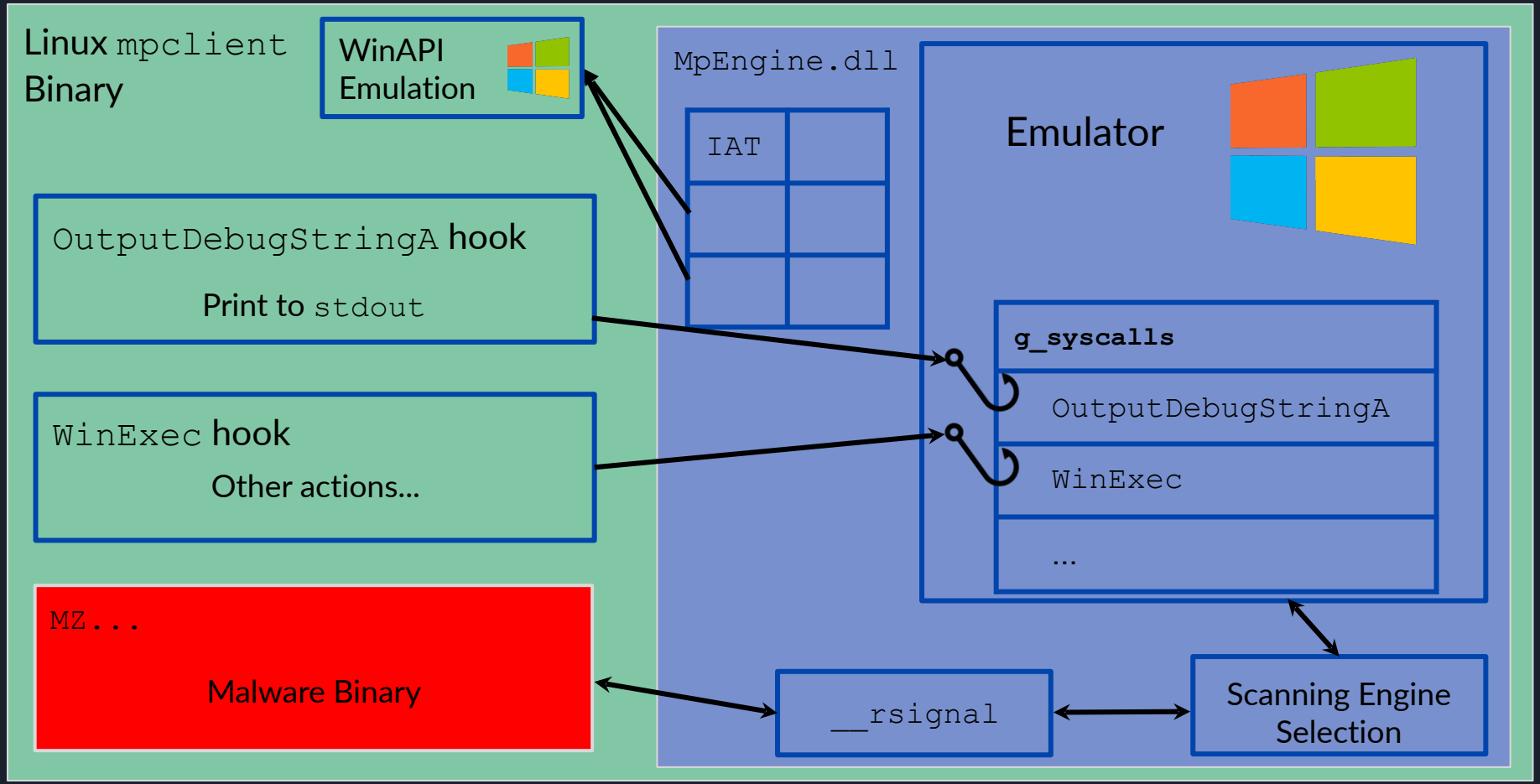
mpclient git.io/fbp0X



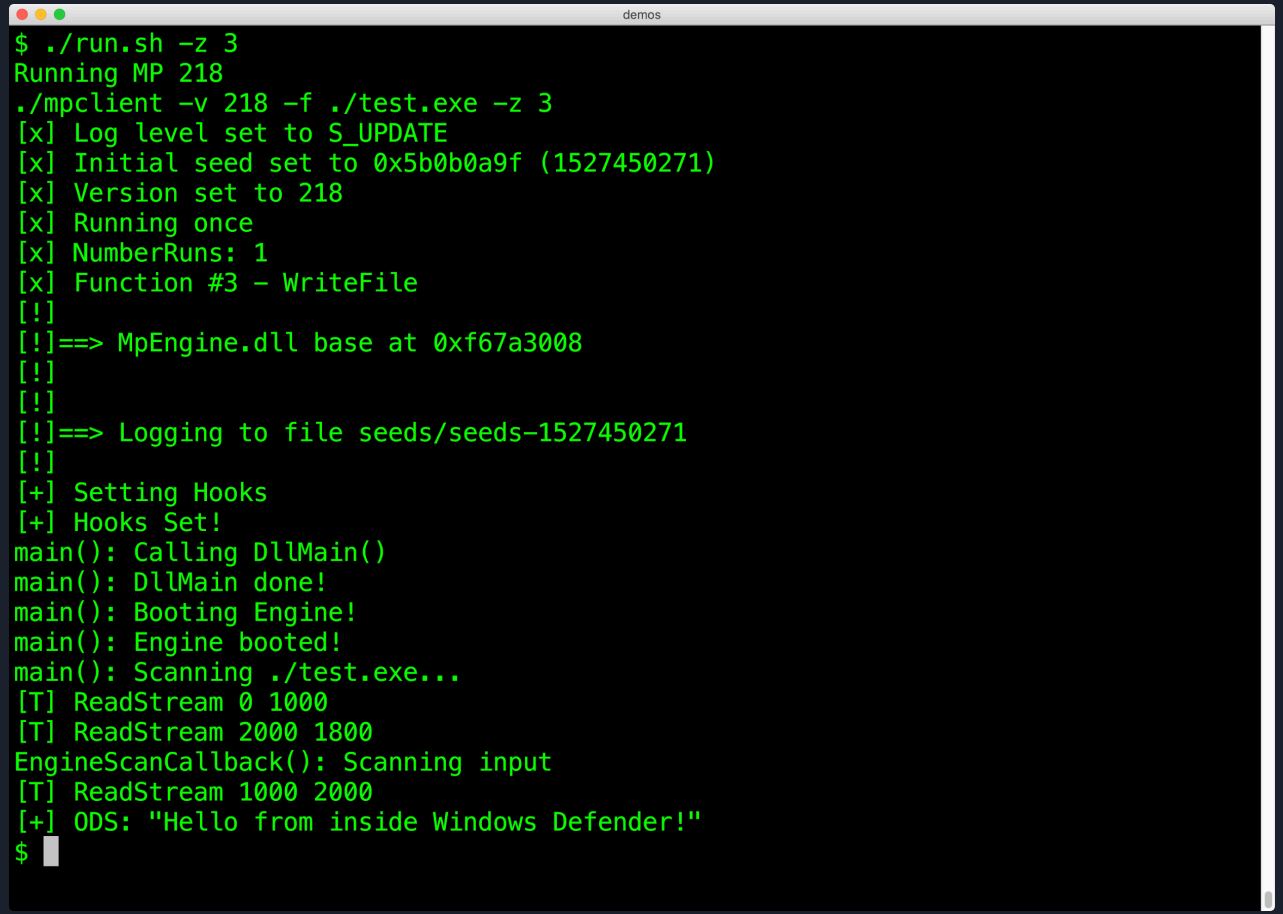
mpclient git.io/fbp0X



Modified `mpclient` - ~3k LoC added `github.com/0xAlexei`



Modified mpclient



`OutputDebugStringA` Hook

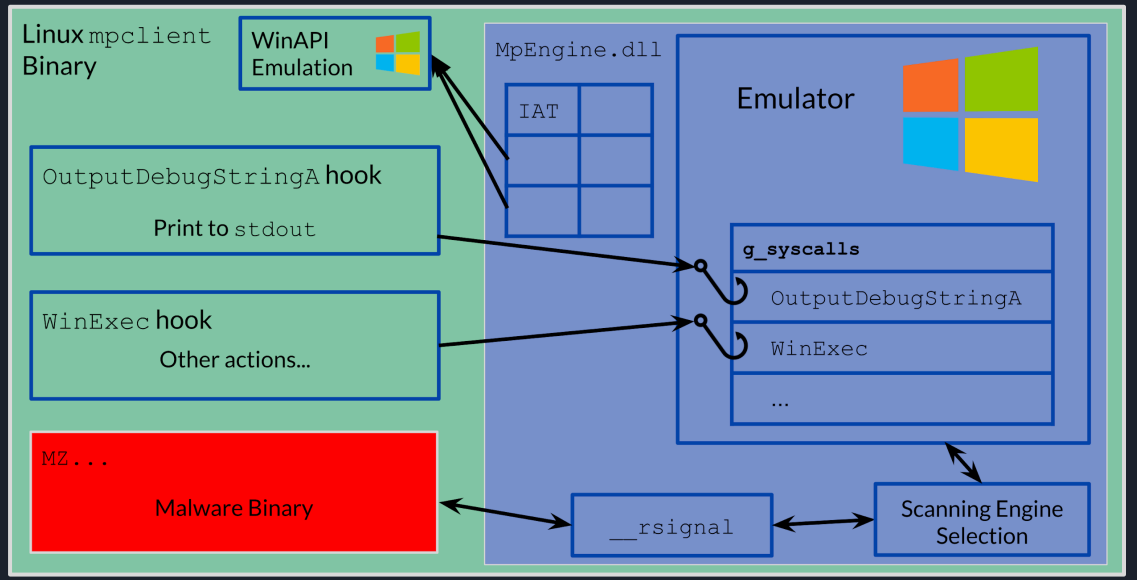
Hook the native function pointer that gets called when `OutputDebugStringA` is called in-emulator

```
void __cdecl KERNEL32_DLL_OutputDebugStringA(pe_vars_t *v)
{
  Parameters<1> arg; // [esp+4h] [ebp-Ch]

  Parameters<1>::Parameters<1>(&arg, v);
  v->m_pDTc->m_vticks64 += 32i64;
}
```

Use existing functions in Defender to interact with function parameters and virtual memory

Mark - Thanks for the idea!

```
RVAS rvas523 = {
    .MPVERNO = "MP_5_23",

    //Parameter functions
    .RVA_Parameters1 = 0x3930f5,
    .RVA_Parameters2 = 0x3b3cfd,
```




```
//OutputDebugString
pOutputDebugStringA = imgRVA(pRVAs->RVA_FP_OutputDebugStringA);
elog(S_DEBUG_VV, "OutputDebugStringA:\t\t0x%06x @ 0x%x", pRVAs->RVA_FP_OutputDebugStringA, *(pOutputDebugStringA));
*pOutputDebugStringA = (uint32_t)KERNEL32_DLL_OutputDebugStringA_hook;
elog(S_DEBUG_VV, "OutputDebugStringA Hooked:\t0x%x", *(pOutputDebugStringA));
```

Dealing With Calling Conventions

When calling `mpengine.dll` functions from `mpclient`: Difficulty of interoperability between MSVC and GCC compiled code

- Possible to massage compiler with `__attribute__` annotations

Easier solution - just hand-write assembly thunks to marshall arguments into the correct format

```
ASM_pe_read_string_ex:
    push ebp
    mov ebp, esp

    mov eax, dword [ebp+0x8]    ;1 - fp
    mov ecx, [ebp+0xc]          ;2

    push dword [ebp+0x18]       ;4
    push dword [ebp+0x14]       ;3 hi
    push dword [ebp+0x10]       ;3

    call eax

    add esp, 0xc
    pop ebp
    ret

ASM__mmap_ex:
    push ebp
    mov ebp, esp

    mov eax, dword [ebp+0x8]; fp
    mov ecx, [ebp+0xc]      ; 2 - v
    mov edx, [ebp+0x10]     ; (SIZE)

    push dword [ebp+0x1c]   ; rights
    push dword [ebp+0x18]   ; addr hi
    push dword [ebp+0x14]   ; addr low

    call eax

    add esp, 0xc
    pop ebp
    ret
```

Dealing With Calling Conventions

When calling `mpengine.dll` functions from `mpclient`: Difficulty of interoperability between MSVC and GCC compiled code

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Easier solution - just hand-write assembly thunks to marshall arguments into the correct format

```
BYTE * __fastcall __mmap_ex
(
    pe_vars_t * v,          // ecx
    unsigned int64 addr,    // too big for edx
    unsigned long  size,    // edx
    unsigned long  rights
);
```

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    mov eax, dword [ebp+0x8]; fp
    mov ecx, [ebp+0xc]      ; 2 - v
    mov edx, [ebp+0x10]     ; (SIZE)

    push dword [ebp+0x1c]   ; rights
    push dword [ebp+0x18]   ; addr hi
    push dword [ebp+0x14]   ; addr low

    call eax

    add esp, 0xc
    pop ebp
    ret
```

Dealing With Calling Conventions

When calling `mpengine.dll` functions from `mpclient`: Difficulty of interoperability between MSVC and GCC compiled code

- Possible to massage compiler with `__attribute__` annotations

Easier solution - just hand-write assembly thunks to marshall arguments into the correct format

```
BYTE * __fastcall __mmap_ex
(
    pe_vars_t * v,          // ecx
    unsigned int64 addr,    // too big for edx
    unsigned long  size,    // edx
    unsigned long  rights
);
```

```
ASM_pe_read_string_ex:
    push ebp
    mov ebp, esp

    mov eax, dword [ebp+0x8]    ;1 - fp
    mov ecx, [ebp+0xc]          ;2

    push dword [ebp+0x18]       ;4
    push dword [ebp+0x14]       ;3 hi
    push dword [ebp+0x10]       ;3

    call eax

    add esp, 0xc
    pop ebp
    ret

ASM__mmap_ex:
    push ebp
    mov ebp, esp

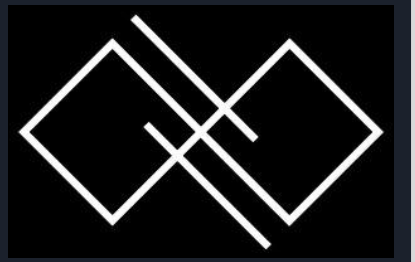
    mov eax, dword [ebp+0x8]; fp
    mov ecx, [ebp+0xc]      ; 2 - v
    mov edx, [ebp+0x10]     ; (SIZE)

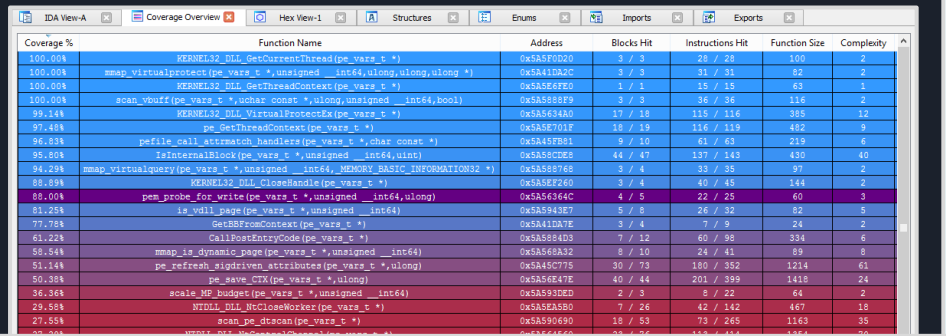
    push dword [ebp+0x1c]   ; rights
                            ; addr hi
                            ; addr low
```

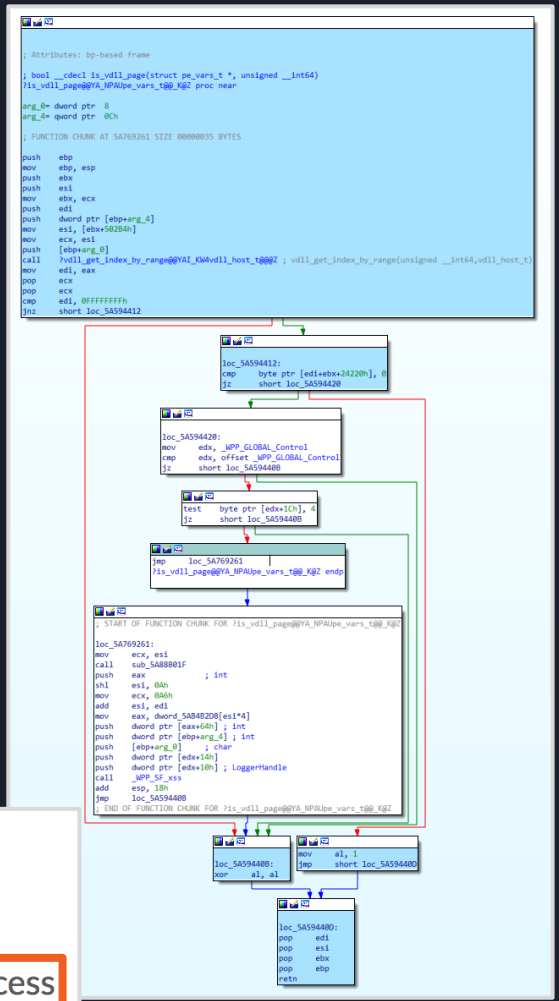
```
// mmap a virtual address
void * e_mmap(void * V, uint64_t Addr, uint32_t Len, uint32_t Rights)
{
    //trampoline through assembly with custom calling convention
    return ASM__mmap_ex(FP__mmap_ex, V, Len, Addr, Rights);
}
```

Dynamic Analysis - Code Coverage

- Getting an overview of what subsystems are being hit is helpful in characterizing a scan or emulation session
  - Breakpoints are too granular
- Emulator has no output other than malware identification
- Lighthouse code coverage plugin for IDA Pro from Markus Gaasedelen of Ret2 Systems / RPISEC







Examples:

Halvar Flake's SSTIC 2018 keynote

- Getting coverage traces from MPENGINE.DLL - difficult because of privileged process

x86_common_context::emulate_CPUID

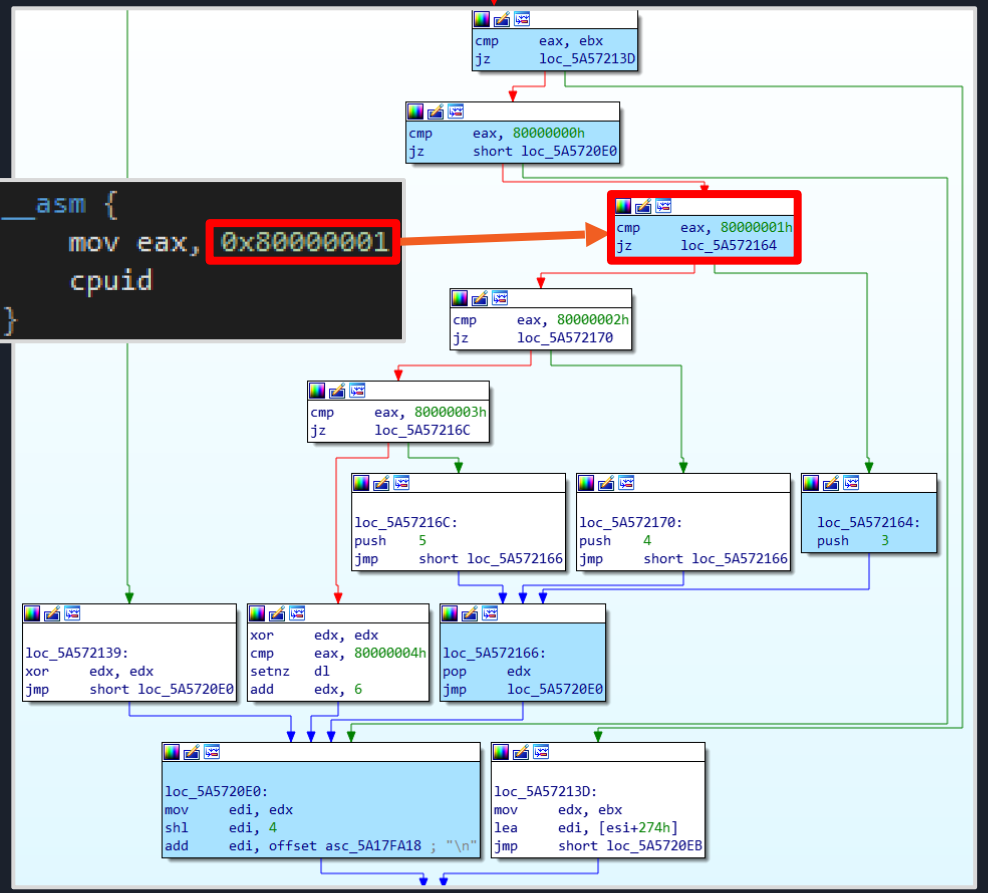
```
; Attributes: bp-based frame

; unsigned int __thiscall x86_common_context::emulate_CPUID(x86_common_context *this, struct DT_context *, bool)
?emulate_CPUID@x86_common_context@@QAEKPAVDT_context@@_N@Z proc near

var_4= dword ptr -4
arg_0= dword ptr  8
arg_4= byte ptr  0Ch

push    ebp
mov     ebp, esp
push    ecx
mov     eax, [ebp+arg_0]
push    ebx
push    esi
mov     esi, ecx
push    edi
push    2
pop     edx
add     dword ptr [esi+3A8h], 100h
mov     ecx, [esi+130h]
adc     dword ptr [esi+3ACh], 0
xor     ebx, ebx
mov     eax, [eax+3668h]
inc     ebx
and     eax, edx
mov     [ebp+var_4], eax
mov     eax, [ecx]
test    eax, eax
jz      loc_5A572139
```

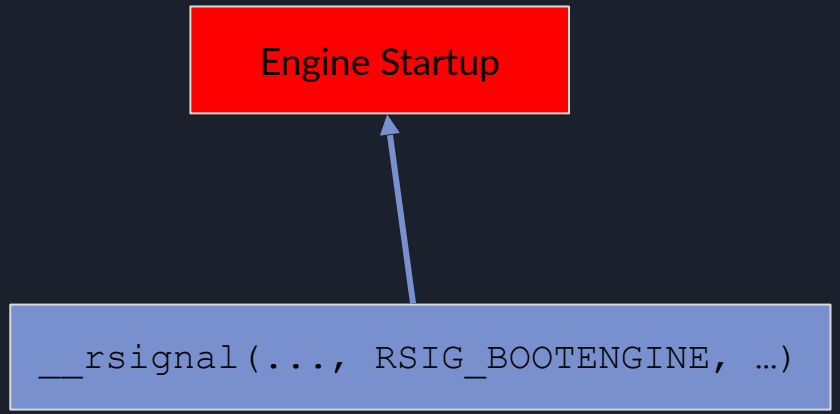
```
__asm {
    mov eax, 0x80000001
    cpuid
}
```



Visualize emulator code coverage when emulating a given “malware” binary

Tracing Timeline

Pintool must be enlightened about custom loaded `mpengine.dll` location - take callback stub ideas from Tavis Ormandy's `deepcover` Pintool
`github.com/taviso/loadlibrary/tree/master/coverage`



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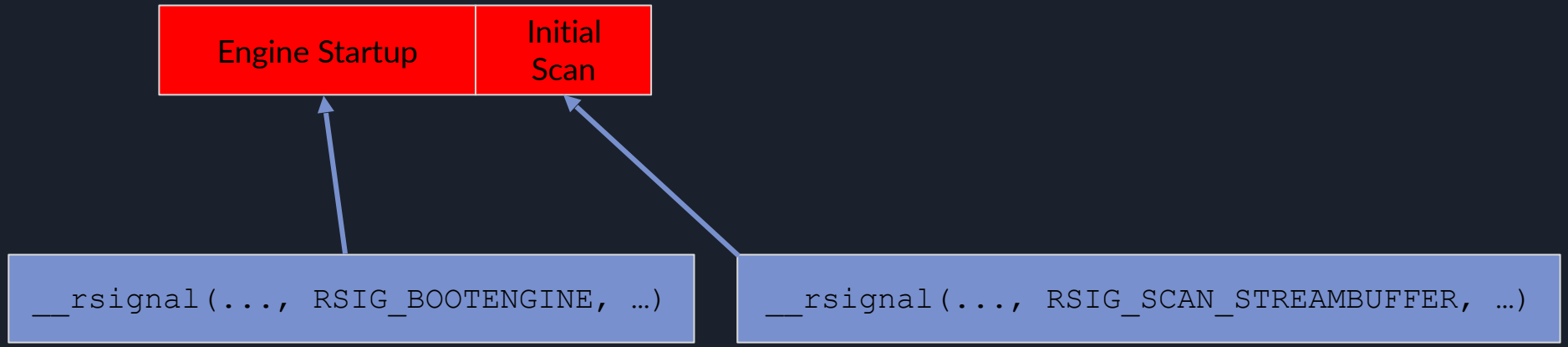

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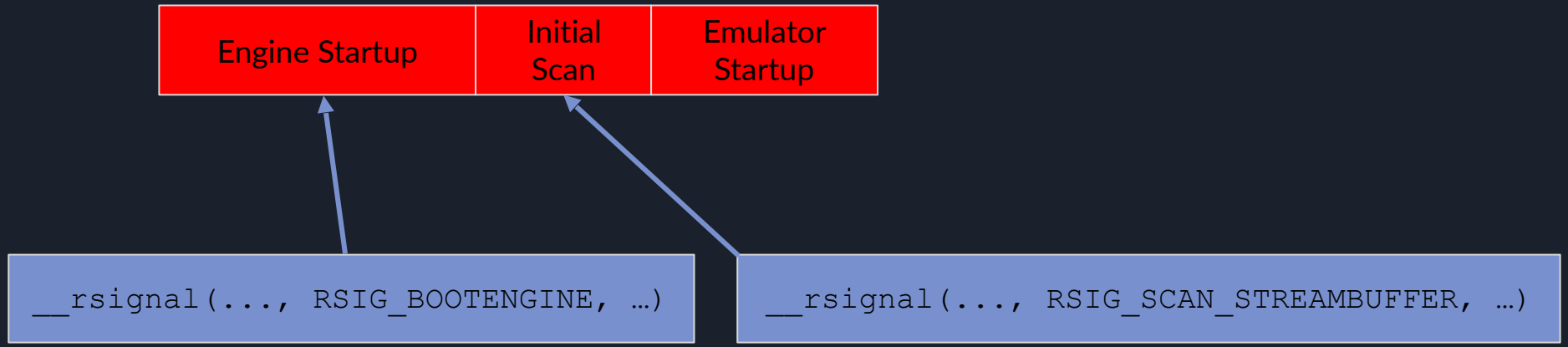

Tracing Timeline

Hooking Defender's emulation functions for `WinExec` and `ExitProcess` allows us to know when emulation starts and stops*

*`ExitProcess` is called at the end of every emulation session automatically - I believe this is because `setup_pe_vstack` puts it at the bottom of the call stack, even for binaries that do not explicitly return to it

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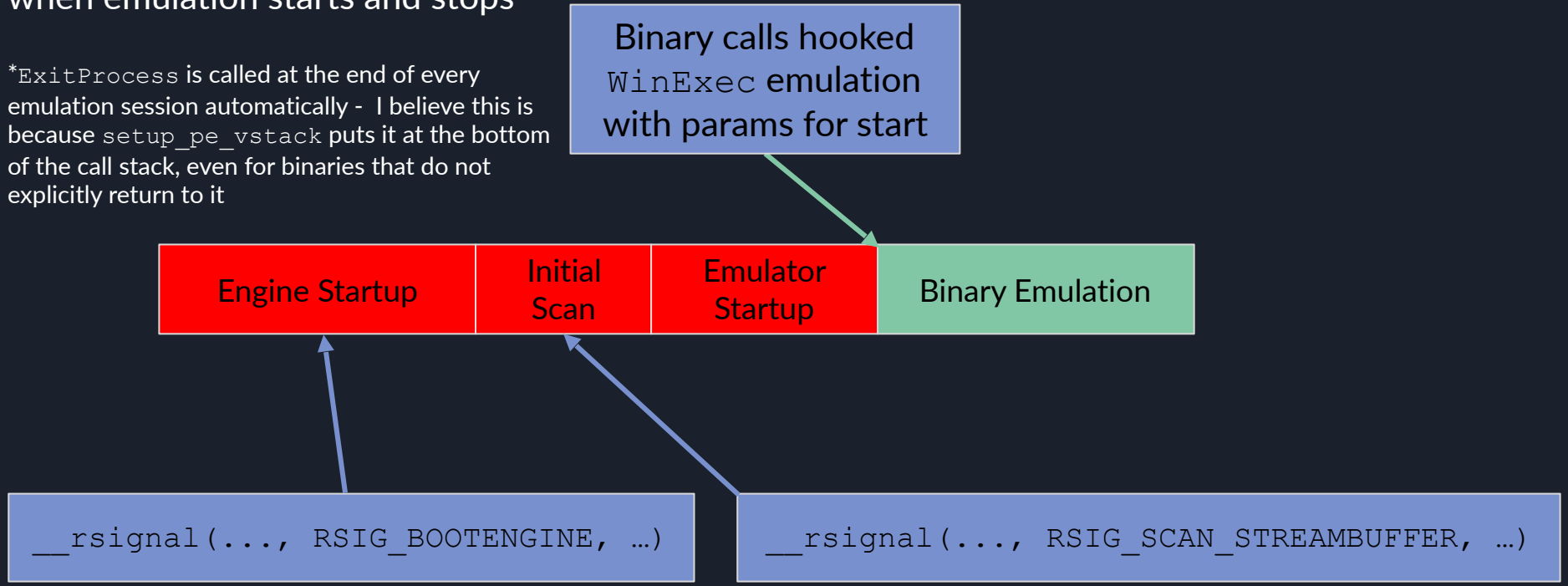

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Hooking Defender's emulation functions for `WinExec` and `ExitProcess` allows us to know when emulation starts and stops*

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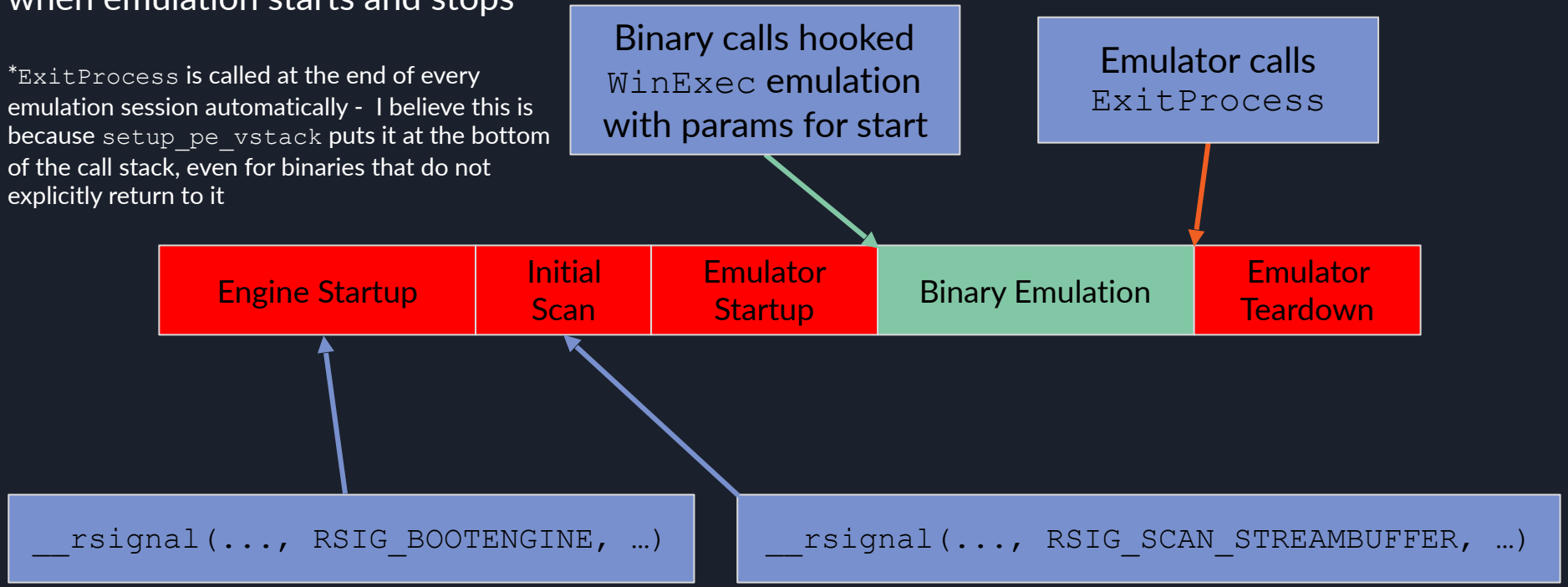

Tracing Timeline

Hooking Defender's emulation functions for `WinExec` and `ExitProcess` allows us to know when emulation starts and stops*

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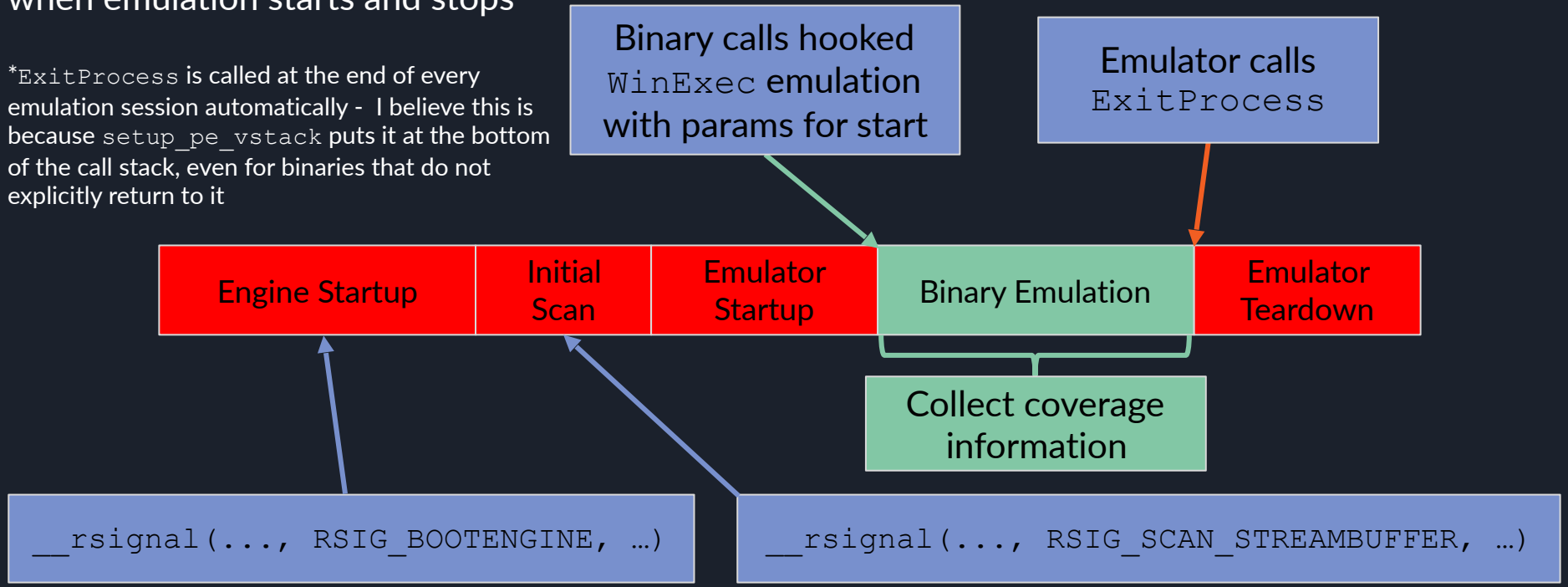


Fuzzing Emulated APIs

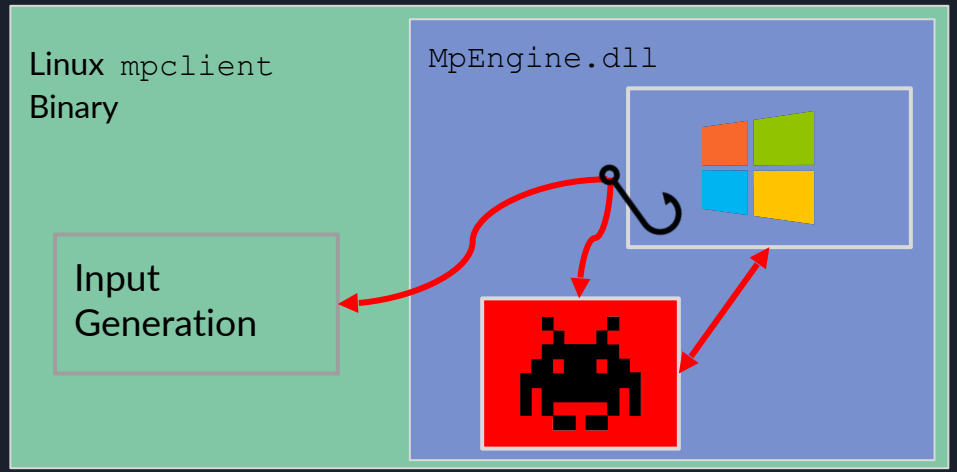
- Create a binary that goes inside the emulator and repeatedly calls hooked `WinExec` function to request new data, then sends that data to functions with native emulations
- Buffers in memory passed to external hook function to populate with parameters
- Could do fuzzing in-emulator too, but this is easier for logging results

```
case ParamTypeDWORD32:
    fuzzParam->Params[i].RawParam = GetFuzzDWORD();
    elog(S_INFO, "\t%d DWORD RawParam: 0x%llx", i, currentParam->RawParam);
    flog(fuzzParam->Init.logfiletest, "\tDWORD: 0x%llx\n", currentParam->RawParam);
    break;

case ParamTypeWORD16:
    fuzzParam->Params[i].RawParam = GetFuzzWORD();
    elog(S_INFO, "\t%d WORD RawParam: 0x%x", i, currentParam->RawParam);
    flog(fuzzParam->Init.logfiletest, "\tWORD: 0x%x\n", currentParam->RawParam);
    break;

case ParamTypeBYTE8:
    fuzzParam->Params[i].RawParam = GetFuzzBYTE();
    elog(S_INFO, "\t%d BYTE RawParam: 0x%x", i, currentParam->RawParam);
    flog(fuzzParam->Init.logfiletest, "\tBYTE: 0x%x\n", currentParam->RawParam);
    break;

case ParamTypeINVALID:
default:
    elog(S_ERROR, "\t%d UNKNOWN 0x%x", i, currentParam->Type);
    fuzzParam->KillSelf = 1;
    break;
```



Input Generation

- Borrow OSX syscall fuzzer code from MWR Labs OSXFuzz project*
- Nothing fancy, just throw random values at native emulation handlers
- Re-seed `rand()` at the start of each emulation session, just save off seeds in a log

*github.com/mwrlabs/OSXFuzz

```
uint32_t GetFuzzDWORD()
{
    int32_t n = 0;

    switch (rand() % 10) {
        case 0:
            switch (rand() % 11)
        {
            case 0:
                n = 0x80000000 >> (rand() & 0x1f);     // 2^n (1 -> 0x10000)
                break;
            case 1:
                n = rand();                            // 0 -> RAND_MAX (likely 0x7fffffff)
                break;
            case 2:
                n = (unsigned int)0xff << (4 * (rand() % 7));
                break;
            case 3:
                n = 0xffff0000;
                break;
            case 4:
                n = 0xfffe000;
                break;
            case 5:
                n = 0xffffff00 | (rand() & 0xff);
                break;
            case 6:
                n = 0xffffffff - 0x1000;
                break;
            case 7:
                n = 0x1000;
                break;
            case 8:
                n = 0x1000 * ((rand() % (0xffffffff / 0x1000)) + 1);
                break;
            case 9:
                n = 0xffffffff;                        // max
                break;
            case 10:
                n = 0x7fffffff;
                break;
        }
```

`NtWriteFile` Overflow

`NtWriteFile` is normally accessible and exported by `ntdll.dll`

- `VFS_Write` has to be triggered with special `apicall`

Tavis' inputs get sanitized out by `NtWriteFileWorker` before it calls down to `VFS_Write`

```
byteOffsLow = 0;
byteOffsHigh = v16->vfptr[1].postDecOpenCount(&v16->vfptr);
hFile = (v16->vfptr[1].__vecDelDtor)(v16);
if ( !VFS_Write(v->vfs, hFile, pBuffer, arg.m_Arg[6].val32, byteOffsHigh, &byteOffsLow) || !byteOffsLow )
  goto LABEL_31;
```

```
LARGE_INTEGER L;
L.QuadPart =
0x2ff9ad29fffffc25;

NtWriteFile(
     hFile,
          NULL,
          NULL,
          NULL,
          &ioStatus,
          buf,
          0x1,
          &L,
          NULL);

L.QuadPart = 0x29548af5d7b3b7c;
NtWriteFile(
     hFile,
          NULL,
          NULL,
          NULL,
          &ioStatus,
          buf,
          0x1,
          &L,
          NULL);
```

NtWriteFile Overflow

NtWriteFile is normally accessible and exported by ntdll.dll

- VFS_Write has to be triggered with special apicall

Tavis' inputs get sanitized out by NtWriteFileWorker before it calls down to VFS_Write

I fuzzed NtWriteFile:

- ~7 minutes @ ~8,000 NtWriteFile calls / second
- Fuzzed Length arguments
- Reproduced Tavis' crash, alternate easier to reach code path through NtWriteFile

Unfortunately, patches for VFS_Write bug also fixed this

```
byteOffsLow = 0;
byteOffsHigh = v16->vfptr[1].postDecOpenCount(&v16->vfptr);
hFile = (v16->vfptr[1].__vecDelDtor)(v16);
if ( !VFS_Write(v->vfs, hFile, pBuffer, arg.m_Arg[6].val32, byteOffsHigh, &byteOffsLow) || !byteOffsLow )
  goto LABEL_31;
```

```
LARGE_INTEGER L;
L.QuadPart =
0x2ff9ad29fffffc25;

NtWriteFile(
     hFile,
          NULL,
          NULL,
          NULL,
          &ioStatus,
          buf,
          0x1,
          &L,
          NULL);

L.QuadPart = 0x29548af5d7b3b7c;
NtWriteFile(
     hFile,
          NULL,
          NULL,
          NULL,
          &ioStatus,
          buf,
          0x1,
          &L,
          NULL);
```

apicall

Custom "apicall" opcode used to trigger native emulation routines

```
0F FF F0 [4 byte immediate]
```

apicall instructions can be disassembled with an IDA Processor Extension Module

```
                        apicall_kernel32_OutputDebugStringA proc near
                                                                ; CODE XREF
8B FF                                   mov     edi, edi
E8 00 00 00 00                          call    $+5
83 C4 04                                add     esp, 4
0F FF F0 BB 14 80 B2                    apicall kernel32!OutputDebugStringA
C2 04 00                                retn    4
                        apicall_kernel32_OutputDebugStringA endp
```

apicall

Custom "apicall" opcode used to trigger native emulation routines

```
0F FF F0 [4 byte immediate]
```

```
immediate = crc32(DLL name, all caps) ^ crc32(function name)
```

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                        apicall_kernel32_OutputDebugStringA proc near
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0F FF F0 BB 14 80 B2            apicall kernel32!OutputDebugStringA
C2 04 00                        retn    4
                        apicall_kernel32_OutputDebugStringA endp
```

apicall

Custom “apicall” opcode used to trigger native emulation routines

```
0F FF F0 [4 byte immediate]
```

```
immediate = crc32(DLL name, all caps) ^ crc32(function name)
```

```
0xB28014BB = crc32(“KERNEL32.DLL”) ^ crc32(“OutputDebugStringA”)
```

apicall instructions can be disassembled with an IDA Processor Extension Module

```
                        apicall_kernel32_OutputDebugStringA proc near
                                                                ; CODE XREF
8B FF                           mov     edi, edi
E8 00 00 00 00                  call    $+5
83 C4 04                        add     esp, 4
0F FF F0 BB 14 80 B2            apicall kernel32!OutputDebugStringA
C2 04 00                        retn    4
                        apicall_kernel32_OutputDebugStringA endp
```

apicall

Custom “apicall” opcode used to trigger native emulation routines

```
0F FF F0 [4 byte immediate]
```

```
immediate = crc32(DLL name, all caps) ^ crc32(function name)
```

```
0xB28014BB = crc32(“KERNEL32.DLL”) ^ crc32(“OutputDebugStringA”)
```

```
0F FF F0 BB 14 80 B2
apicall kernel32!OutPutDebugStringA
```

apicall instructions can be disassembled with an IDA Processor Extension Module

```
                        apicall_kernel32_OutputDebugStringA proc near
                                                                ; CODE XREF
8B FF                           mov     edi, edi
E8 00 00 00 00                  call    $+5
83 C4 04                        add     esp, 4
0F FF F0 BB 14 80 B2            apicall kernel32!OutputDebugStringA
C2 04 00                        retn    4
                        apicall_kernel32_OutputDebugStringA endp
```

Locking Down `apicall`

`is_vdll_page` call added to `__call_api_by_crc` in 6/20/2017 `mpengine.dll` build - is the `apicall` instruction coming from a VDLL?

```
aX64                db '{x64}',0                ; DATA XR
                    align 4
aPea_invalid_ap     db 'pea_invalid_apicall_opcode' 0
                                                ; DATA XR
                    align 4
aKernel32_dll_0     db 'kernel32.dll',0         ; DATA XR
```

New AV heuristic trait added

Can't just trigger `apicall` from malware `.text` section or otherwise malware-created memory (eg: rwx allocation) anymore

```
if ( !*(v_pe_vars + 167453) )
{
  LODWORD(page) = v6;
  if ( is_vdll_page(v_alias, page) && (!mmap_is_dynamic_page(v_alias, *(&v26 - 1)) || nidsearchrecid(v29) != 1) )
  {
    if ( !*(v_pe_vars + 167454) )
    {
      qmemcpy(&dst, &NullSha1, 0x14u);
      v15 = *v_pe_vars;
      MpSetAttribute(0, 0, &dst, 0, *(&v27 - 1));
      *(v_pe_vars + 167454) = 1;
    }
    return 0;
  }
}
v16 = &syscall_table;
do
{
  v17 = &v16[2 * (v13 / 2)];
  if ( *(v17 + 4) >= v29 )
  {
```

If `apicall` did not come from a VDLL, set a heuristic and deny it

Proceed with processing if `apicall` is ok

Bypass

- `apicall` stubs are located throughout VDLLs
- They can be located in memory and called directly by malware, with attacker controlled arguments
  - Passes `is_vdll_page` checks

Response from MSFT: “We did indeed make some changes to make this interface harder to reach from the code we’re emulating -however, that was never intended to be a trust boundary.

Accessing the internal APIs exposed to the emulation code is not a security vulnerability…”

```
text:7C816E1E 8B FF                                  mov     edi, edi
text:7C816E20 E8 00 00 00 00                         call    $+5
text:7C816E25 83 C4 04                               add     esp, 4
text:7C816E28 0F FF F0 3C 28 D6 CC                   apicall ntdll!VFS_SetLength
text:7C816E2F C2 08 00                               retn    8
text:7C816E32                          ; ---------------------------------------------------------------------------
text:7C816E32 8B FF                                  mov     edi, edi
text:7C816E34 E8 00 00 00 00                         call    $+5
text:7C816E39 83 C4 04                               add     esp, 4
text:7C816E3C 0F FF F0 41 3B FA 3D                   apicall ntdll!VFS_GetLength
text:7C816E43 C2 08 00                               retn    8
text:7C816E46                          ; ---------------------------------------------------------------------------
text:7C816E46 8B FF                                  mov     edi, edi
text:7C816E48 E8 00 00 00 00                         call    $+5
text:7C816E4D 83 C4 04                               add     esp, 4
text:7C816E50 0F FF F0 FC 99 F8 98                   apicall ntdll!VFS_Read
text:7C816E57 C2 14 00                               retn    14h
text:7C816E5A                          ; ---------------------------------------------------------------------------
text:7C816E5A 8B FF                                  mov     edi, edi
text:7C816E5C E8 00 00 00 00                         call    $+5
text:7C816E61 83 C4 04                               add     esp, 4
text:7C816E64 0F FF F0 E7 E3 EE FD                   apicall ntdll!VFS_Write
text:7C816E6B C2 14 00                               retn    14h
text:7C816E6E                          ; ---------------------------------------------------------------------------
text:7C816E6E 8B FF                                  mov     edi, edi
text:7C816E70 E8 00 00 00 00                         call    $+5
text:7C816E75 83 C4 04                               add     esp, 4
text:7C816E78 0F FF F0 1D 86 73 21                   apicall ntdll!VFS_CopyFile
text:7C816E7F C2 08 00                               retn    8
text:7C816E82                          ; ---------------------------------------------------------------------------
text:7C816E82 8B FF                                  mov     edi, edi
text:7C816E84 E8 00 00 00 00                         call    $+5
text:7C816E89 83 C4 04                               add     esp, 4
text:7C816E8C 0F FF F0 B1 0D B0 47                   apicall ntdll!VFS_MoveFile
text:7C816E93 C2 08 00                               retn    8
text:7C816E96                          ; ---------------------------------------------------------------------------
text:7C816E96 8B FF                                  mov     edi, edi
text:7C816E98 E8 00 00 00 00                         call    $+5
text:7C816E9D 83 C4 04                               add     esp, 4
text:7C816EA0 0F FF F0 4A BD 6E C0                   apicall ntdll!VFS_DeleteFile
text:7C816EA7 C2 04 00                               retn    4
```

Bypass Example

`OutputDebugStringA` can be normally hit from `kernel32`, so this is ultimately just a unique way of doing that

```
VOID OutputDebugStringA_APICALL(PCHAR msg)
{
    typedef VOID(*PODS)(PCHAR);
    HMODULE k32base = LoadLibraryA("kernel32.dll");
    PODS apicallODS = (PODS)((PBYTE)k32base + 0x16d4e);
    apicallODS(msg);
}
```

`Kernel32` base offset: `0x16d4e`

Comes from `kernel32` VDLL, so passes `is_vdll_page` checks

```
apicall_kernel32_OutputDebugStringA proc near
                                        ; CODE XREF:
                mov     edi, edi
                call    $+5
                add     esp, 4
                apicall kernel32!OutputDebugStringA
                retn    4
apicall_kernel32_OutputDebugStringA endp
```

Outline

1. Introduction
2. Tooling & Process
3. Discussion
4. Conclusion

Reverse Engineer Intuitions

- It's easy to detect for emulator (or file format unpacker) presence - test an `EICAR` dropper
- Everyone has to emulate `Sleep()` with custom code
- Everyone emulates `cpuid`
- Everyone emulates `rstsc`, but messes up `rdtscp`
- Emulators have lots of strings - these can be found in memory dumps to help identify emulator code

- Everyone builds custom tools when doing offensive research, but this is especially true for AV RE

Reverse Engineer Intuitions - Rolf Rolles in 2013

I've done this same exercise with anti-virus engines on a number of occasions. Generally the steps I use are:

1. Identify the CPU/Windows emulator. This is generally the hardest part. Look at filenames, and also grep the disassembly for large switch statements. Find the switches that have 200 or more cases and examine them individually. At least one of them will be related to decoding the single-byte X86 opcodes.
2. Find the dispatcher for the CALL instruction. Usually it has special processing to determine whether a fixed address is being called. If this approach yields no fruit, look at the strings in the surrounding modules to see anything that is obviously related to some Windows API.
3. Game over. AV engines differ from the real processor and a genuine copy of Windows in many easily-discernible ways. Things to inspect: pass bogus arguments to the APIs and see if they handle erroneous conditions correctly (they never do). See if your emulator models the AF flag. Look up the exception behavior of a complex instruction and see if your emulator implements it properly. Look at the implementations of GetTickCount and GetLastError specifically as these are usually miserably broken.

share improve this answer

answered Sep 18 '13 at 8:00

Rolf Rolles

4,248 ● 17 ● 28

Programmer "Easter Eggs"

aaa_TouchMeNot_.txt — UNREGISTERED

```
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```

```c
signed int __stdcall GetComputerNameExA(signed int NameType, LPCSTR lpBuffer, LPDWORD lpnSize)
{
  if ( NameType >= ComputerNameMax )
  {
    SetError(ERROR_INVALID_PARAMETER);
    return 0;
  }
  if ( !lpnSize || !lpBuffer && *lpnSize )
  {
    SetError(ERROR_INVALID_PARAMETER);
    return 0;
  }
  if ( !NameType
    || NameType == ComputerNameDnsHostname
    || NameType == ComputerNamePhysicalNetBIOS
    || NameType == ComputerNamePhysicalDnsHostname
  {
    if ( *lpnSize < ComputerNameMax )
    {
      *lpnSize = ComputerNameMax;
      SetError(ERROR_MORE_DATA);
      return 0;
    }
    memcpy(lpBuffer, "HAL9TH", 7);
    *lpnSize = 7;
  }
  return 1;
}
```

```javascript
var num = new Number(1);
var node = document.createTextNode("node");
var elem = document.createElement("element");
num.appendChild = elem.appendChild;
num.appendChild(node);

triggerEvent(): err_typeerror
triggerEvent(): error_tostring
Log(): uncaught exception: TypeError: node.insertBefore()
          'this' object must be DOM Object (BUG, should never
happen)
```

In-Emulator Signaling

```
mov  edi, edi          ; WinAPI hot patch point
push ebp               ; function prologue
mov  ebp, esp          ; function prologue
nop
lock mov ebx, 0xff[1b lib #][2b func #]
pop  ebp               ; function epilogue
ret  [size of args]    ; stack cleanup
nop...                 ; nops between functions
```

Figure 7: Example of code extracted from AVG's `kernel32.dll` in memory. The second byte of the `mov` instruction argument denotes the library, while the third and fourth bytes denote a specific function. AVG's CPU emulator presumably intercepts the obscure "`lock mov ebx`", and invokes code to emulate the function.

```
void __stdcall apicall_kernel32_OutputDebugStringA(int a1)
{
  __asm { apicall kernel32!OutputDebugStringA }
}
```

Attackers can discover in-emulator control operations

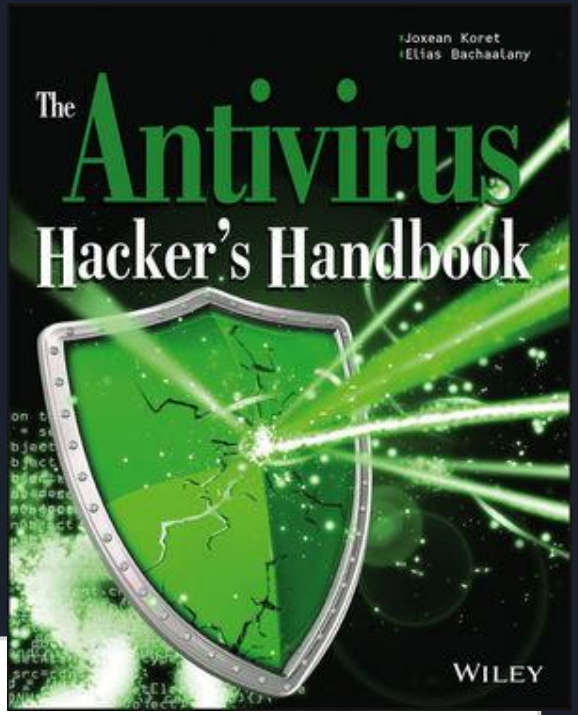
Why not just use `int/syscall`?

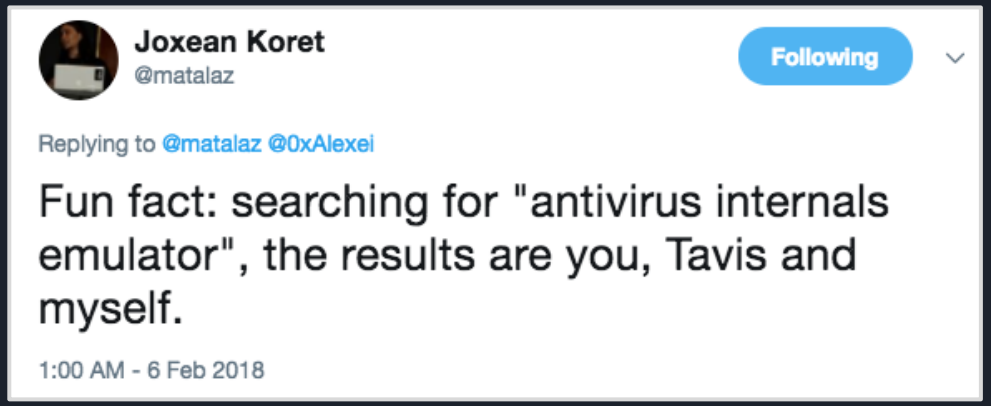
```
; Exported entry  72. CopyFileWWorker


public CopyFileWWorker
CopyFileWWorker proc near
mov     edi, edi
call    $+5
add     esp, 4
apicall kernel32!CopyFileWWorker
retn    0Ch
CopyFileWWorker endp
```

Antivirus Reverse Engineering

- People constantly talk about what AVs can or can't do, and how/where they are vulnerable
- These claims are mostly backed up by Tavis Ormandy's work at Project Zero and a handful of other conference talks, papers, and blog posts

- I hope we'll see more AV research in the future





Security Through Obscurity?

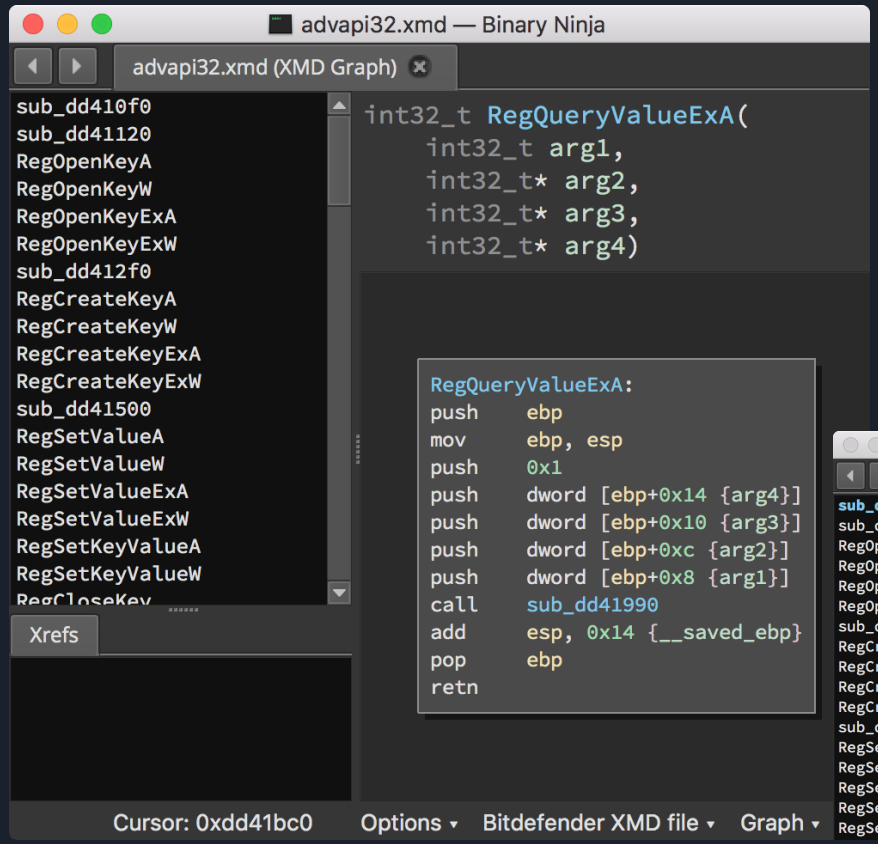
- Preventing reverse engineering is futile
  - Obfuscation and custom binary formats don't stop RE, and can be overcome with one-time effort
  - Side channel analyses like "AVLeak" are also possible

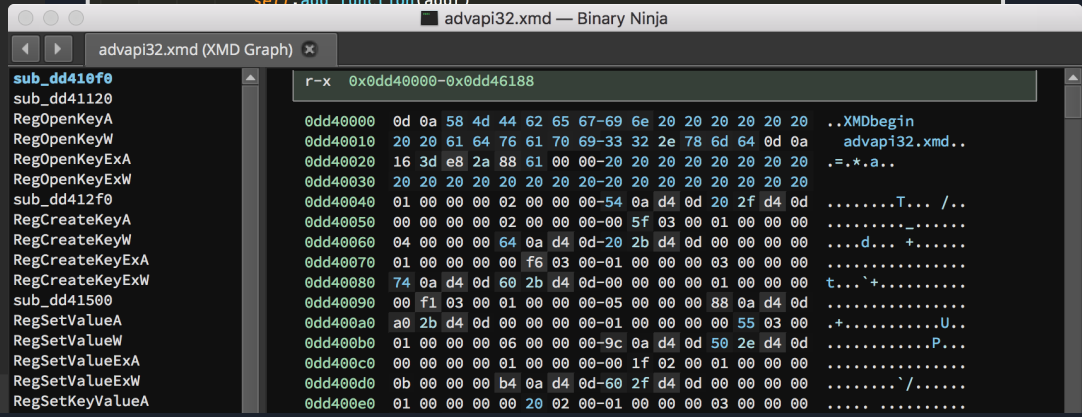
- Introspectibility and debugability are poor → only *motivated competent* adversaries will perform RE
  - Malicious actors *already are* - search any unique string from my presentations - you'll find malware samples from long before I presented

Custom Binary Format Example: Bitdefender XMDs

Custom Binary Ninja loader:
~150 LoC, 4 hours of work

```
def init(self):
    try:
        hdr = self.raw_data.read(0,0x40)
        self.unknown1 = struct.unpack("<I", hdr[0x20:0x24])[0]
        log_info("Unknown 1: " + hex(self.unknown1))
        self.size = struct.unpack("<I", hdr[0x24:0x28])[0]
        log_info("Size: " + hex(self.size))
        self.add_auto_segment(BASE, self.size, 0, self.size,
            SegmentFlag.SegmentReadable|SegmentFlag.SegmentExecutable)

        i = 0
        while True:
            func = self.raw_data.read(0x40 + i*4*7, 7*4)
            args, uk, name, addr, uk2, uk3, uk4 = struct.unpack("<IIIIIII", func)

            if args > BASE and args < BASE + self.size:
                break

            functionname =  self.read(name, 100).split("\x00")[0]
            if addr == 0:
                log_info(functionname + " found, but address is 0")


            #log_error(hex(name) + "'" + functionname + "'")
            else:
                self.add_function(addr)
```





Emulator Exploitation

- Emulators, like web browsers, provide the primitives necessary for modern binary exploitation
- Micro-level: Software attack surface is immense, and the software runs at high privilege on the OS
- Macro-level: For IT organizations, AV software is similar - high privilege within a network, and adds attack surface to your most sensitive assets
- AV engines *seem* intuitively very easy to sandbox

Outline

1. Introduction
2. Tooling & Process
3. Discussion
4. Conclusion

Code & More Information

github.com/0xAlexei

Code release:

- `OutputDebugStringA` hooking
- “Malware” binary to go inside the emulator
- Some IDA scripts, including `apicall` disassembler

Article in PoC||GTFO 0x19:

- `OutputDebugStringA` hooking
- Patch diffing and `apicall` bypass
- `apicall` disassembly with IDA processor extension module

Conclusion

1. I had a great time reverse engineering Windows Defender - seriously cool software
2. REs will create custom tools to address AV complexity
3. Resistance to RE is futile, so be smart about design

JS Engine & Emulator slides:

bit.ly/2qio857

bit.ly/2CxyZ3l

@0xAlexei

Open DMs

Thank You:

- Tavis Ormandy & Natalie Silvanovich @ Google P0 - exposing the engine, `mpclient`, sharing ideas
- Mark - hooking ideas
- Joxean Koret - OG AV hacker
- Virus Bulletin - hosting me and editing my paper

github.com/0xAlexei